

Data-driven Invariant Learning for Probabilistic Programs

Jialu Bao

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Collaborated work with Nitesh Trivedi, Drashti Pathak, Justin Hsu, Subhajit Roy

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Can we automatically find
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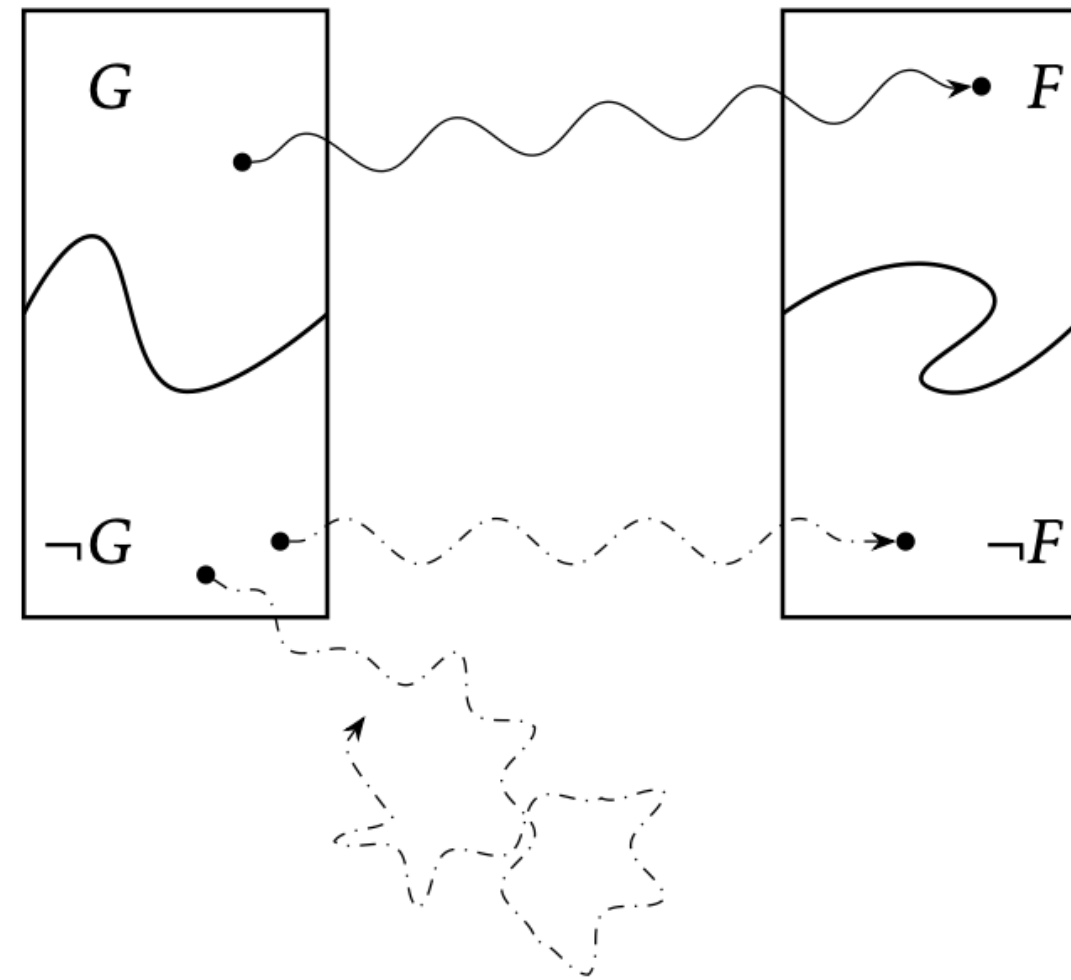
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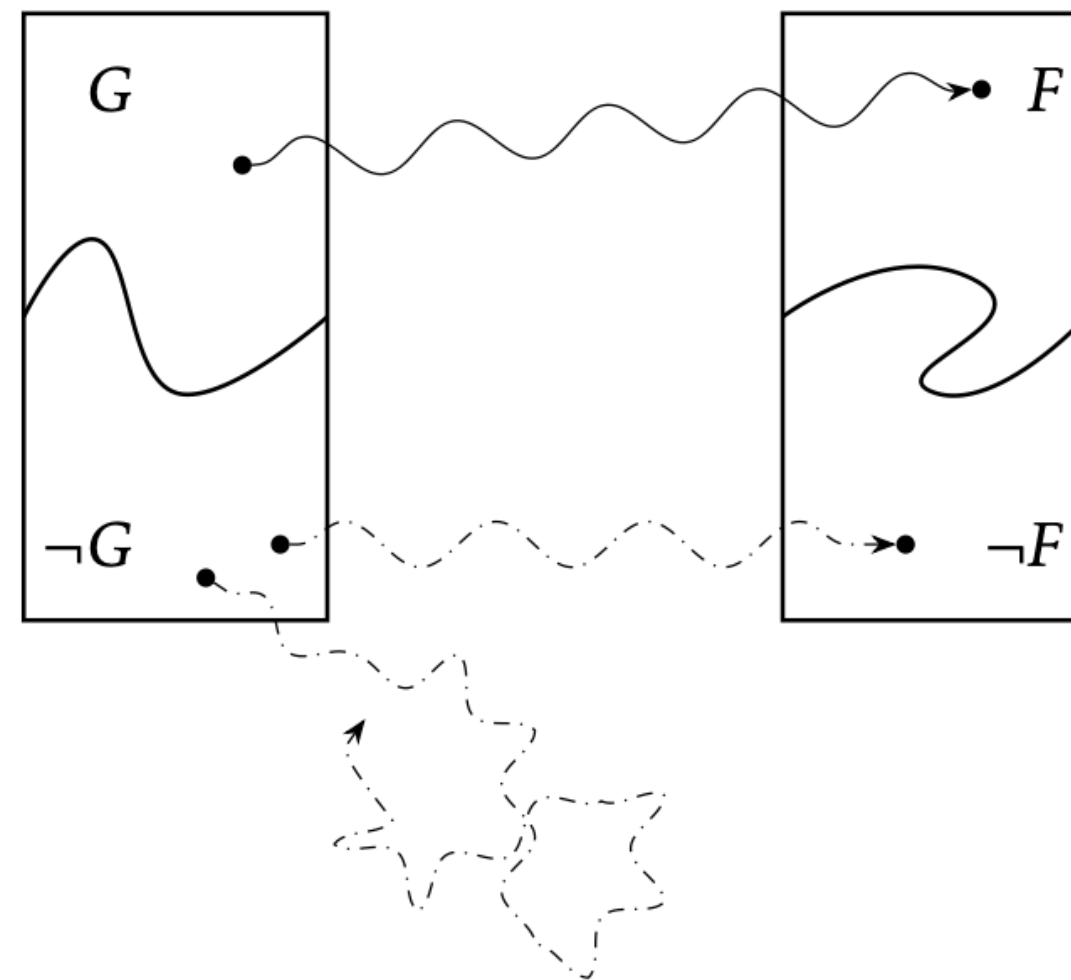
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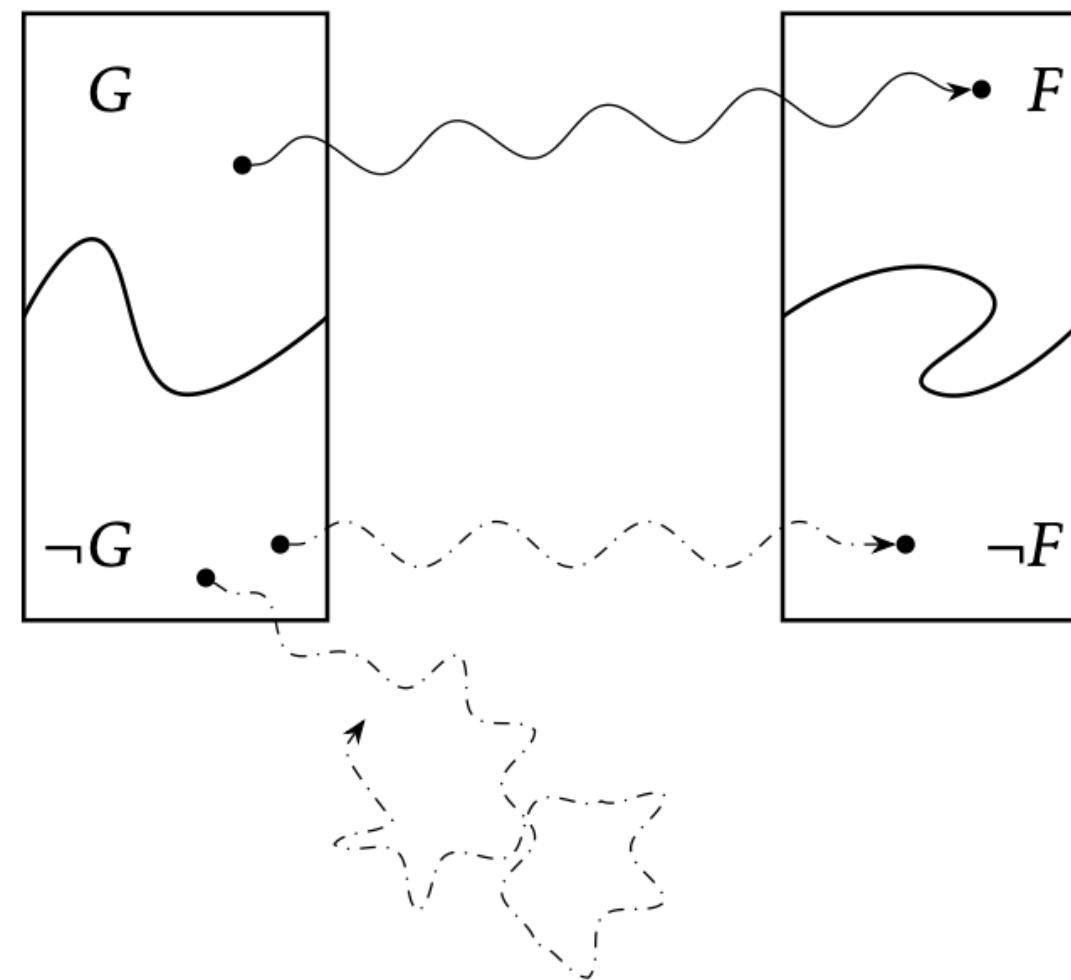
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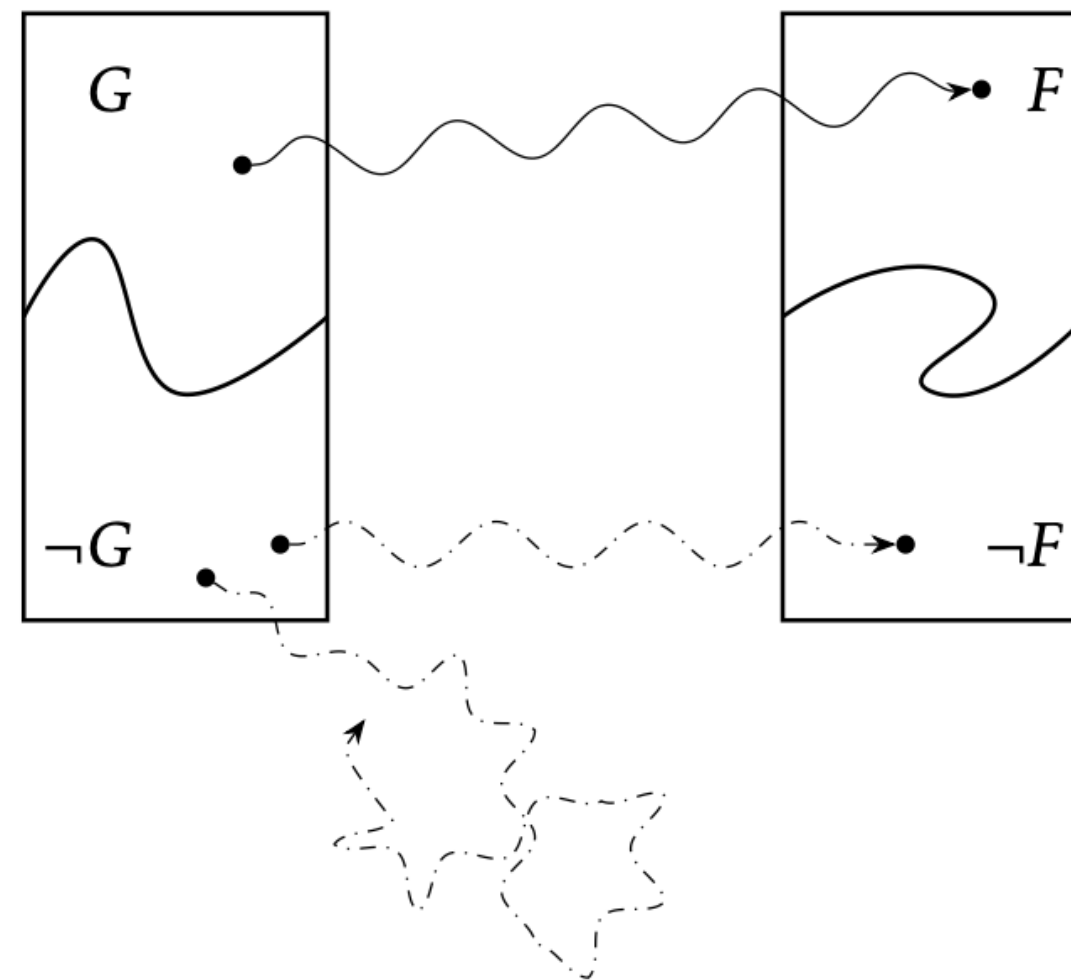


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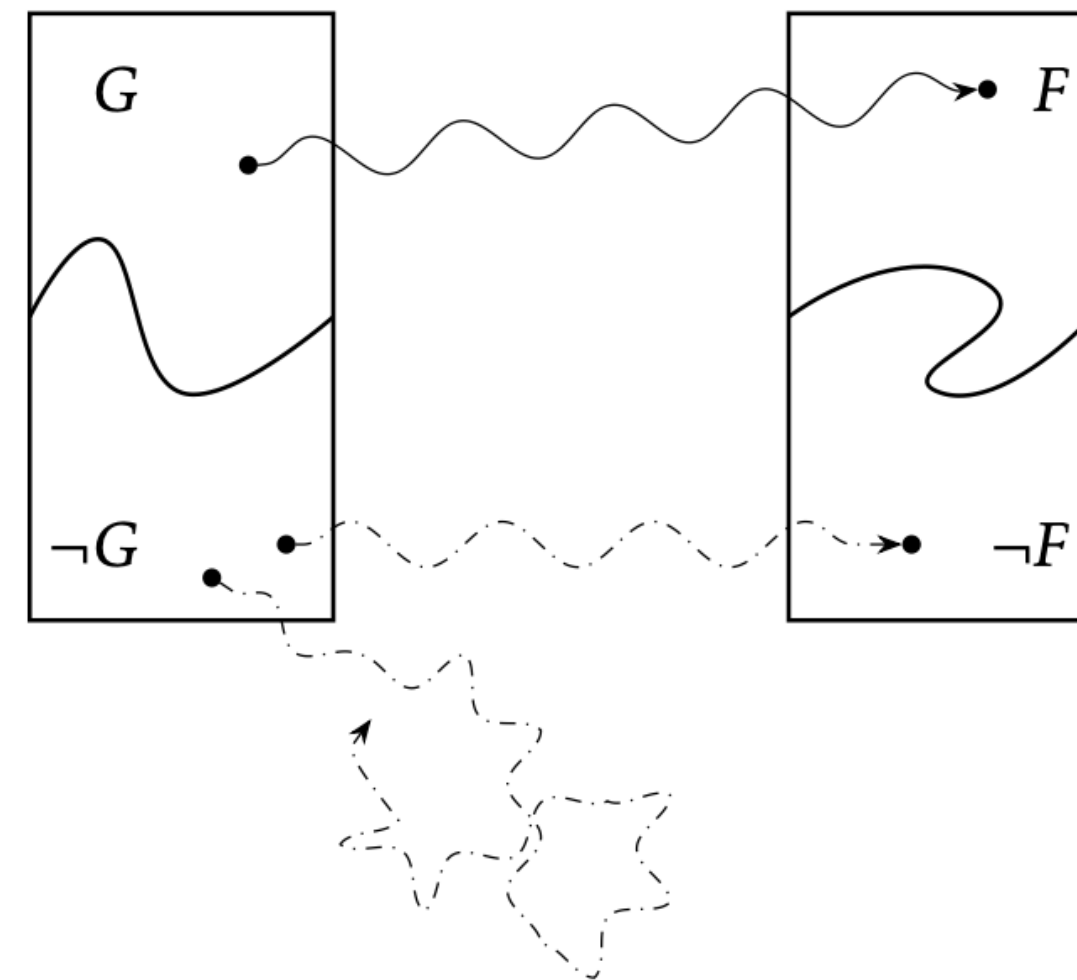
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- Sequencing: $wpc(P; Q, F) := wpc(P, wpc(Q, F))$

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 - Iverson bracket: $[G]$ maps states where the assertion G holds to 1 and maps other states to 0

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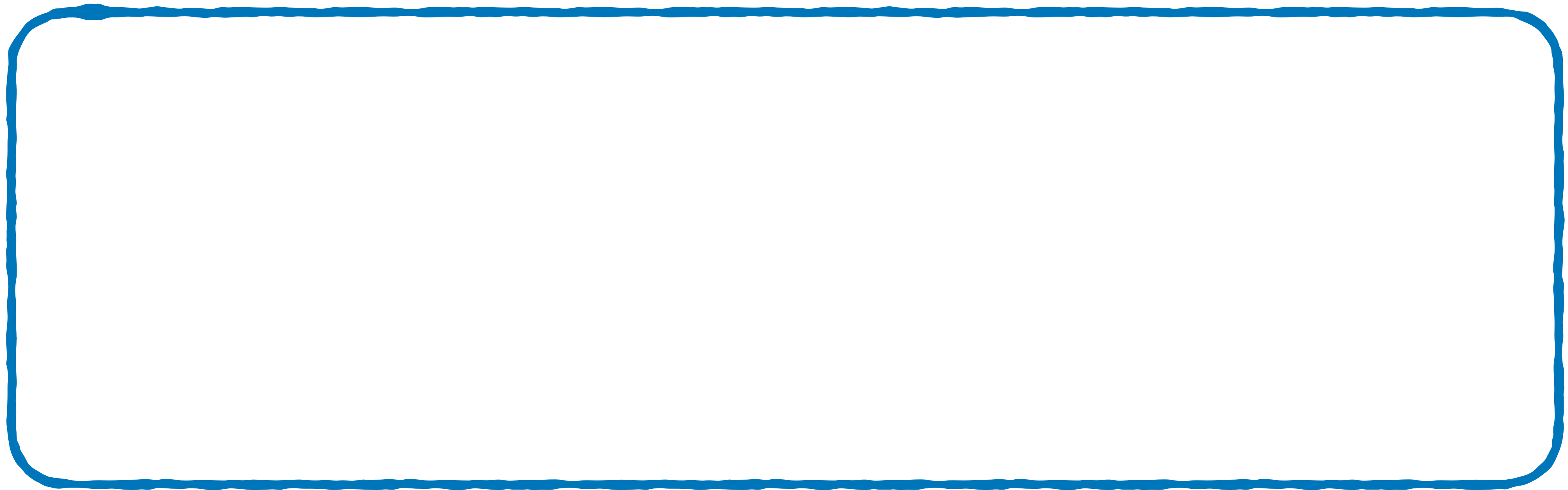
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 - Let the expectation e be $[Ev]$.

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Theorem. $wpe(C, e) = \lambda s. \text{ expected value of } e \text{ after running } C \text{ from } s$

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- Goal: find expectation I such that $I = [G] \cdot wpe(P, I) + [\neg G] \cdot e$.
- We call e the postexpectation, and call I an **exact invariant** of the loop.

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- In other words, in simple cases, I is an exact invariant iff $I = wpe(\text{while } G \text{ do } P, e)$.
- Simple: the loop is almost surely terminating and e is upper bounded by a constant



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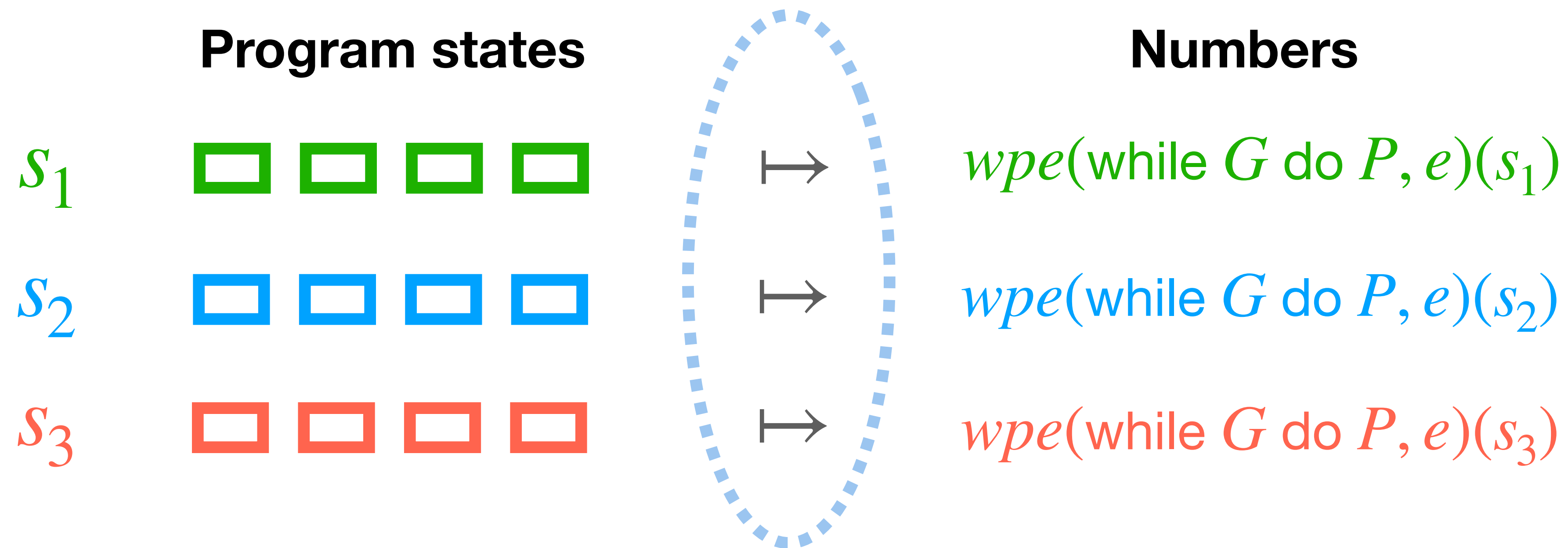
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	Program states		Numbers
s_1		\mapsto	$wpe(\text{while } G \text{ do } P, e)(s_1)$
s_2		\mapsto	$wpe(\text{while } G \text{ do } P, e)(s_2)$
s_3		\mapsto	$wpe(\text{while } G \text{ do } P, e)(s_3)$

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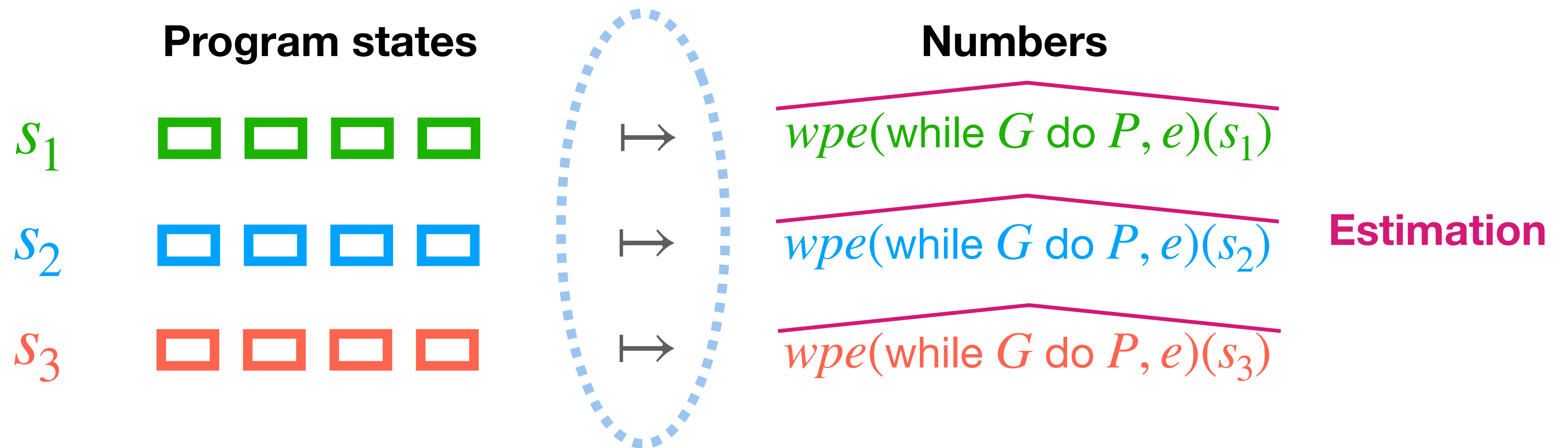
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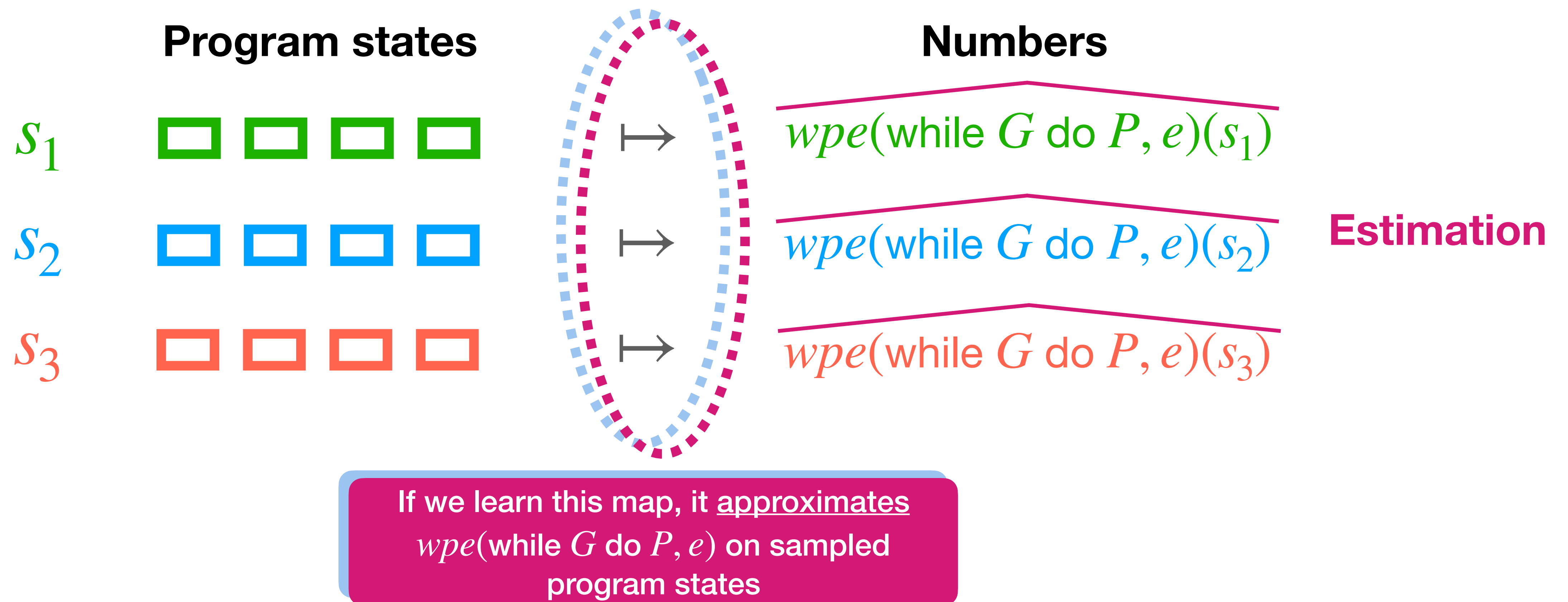
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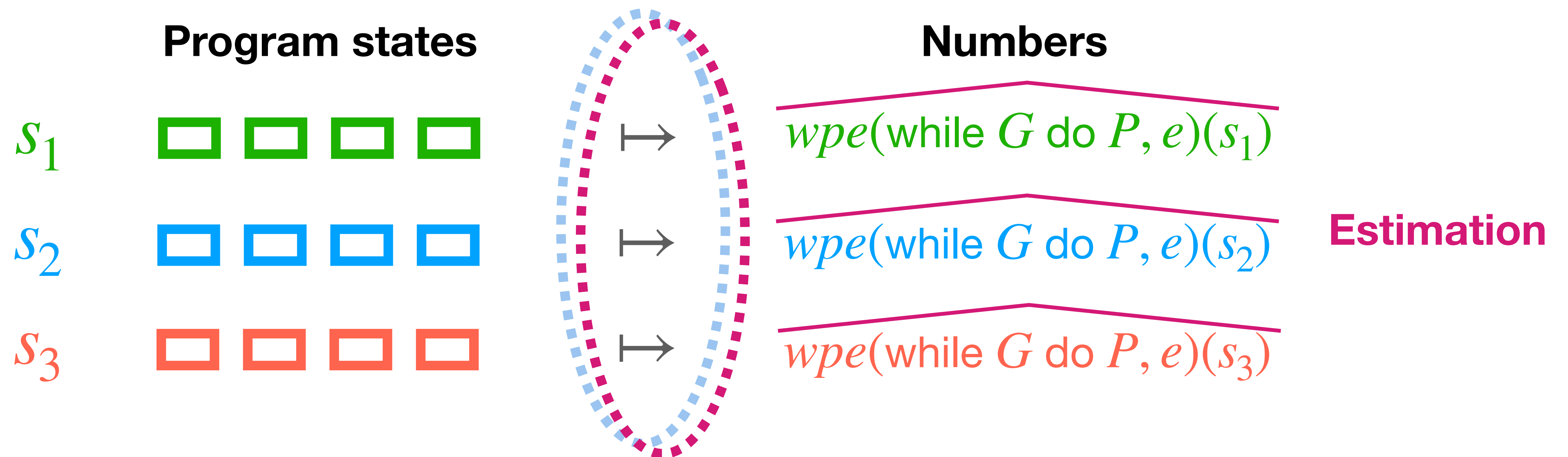
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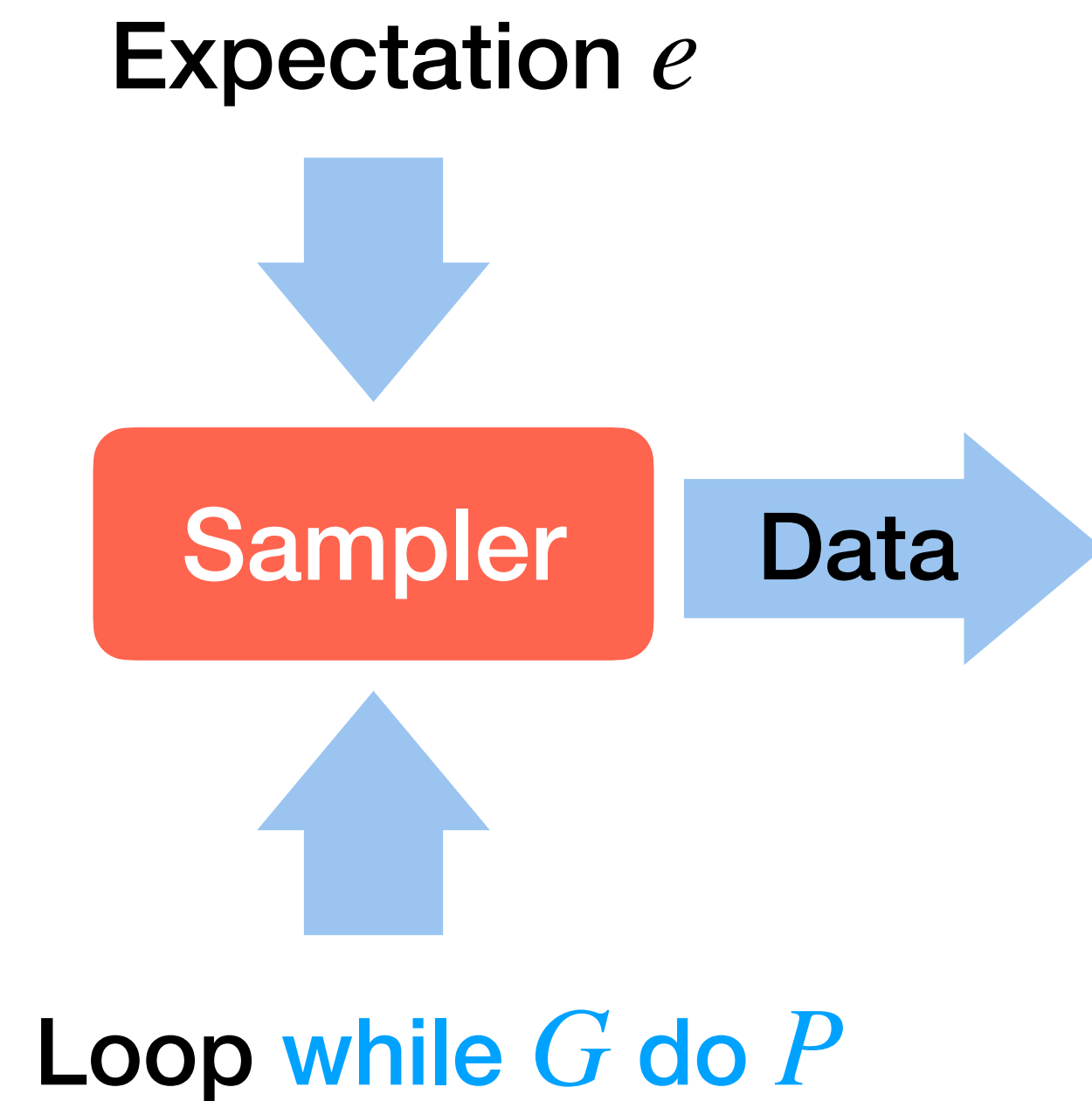
If we learn this map, it approximates $wpe(\text{while } G \text{ do } P, e)$ on sampled program states

The learned map may not be $wpe(\text{while } G \text{ do } P, e)$ but we can check whether it is an exact invariant.

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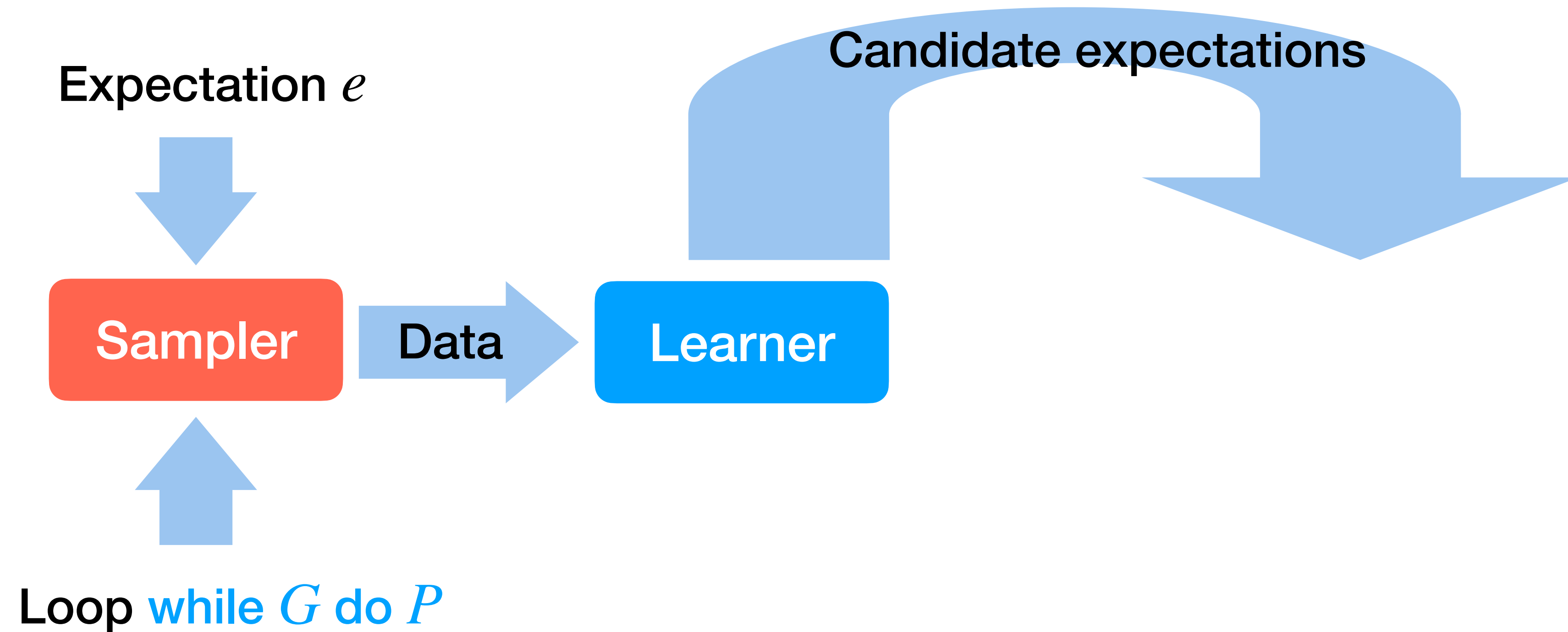
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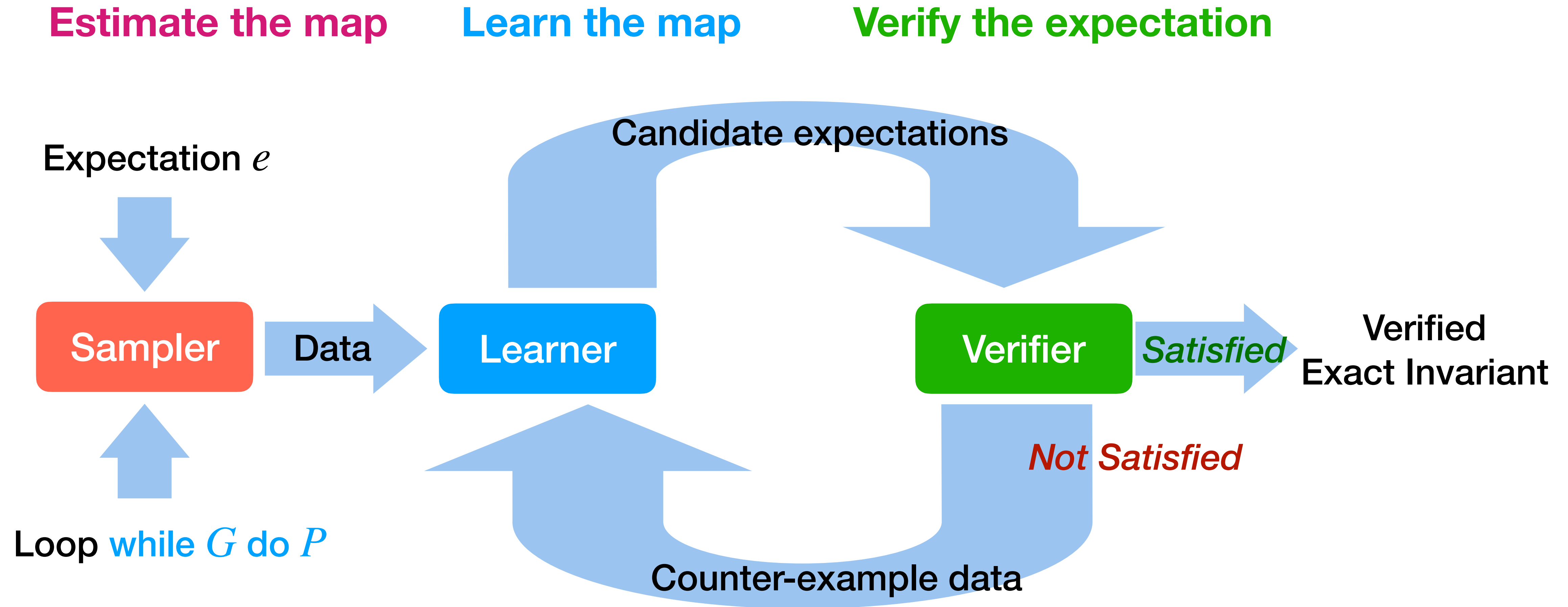
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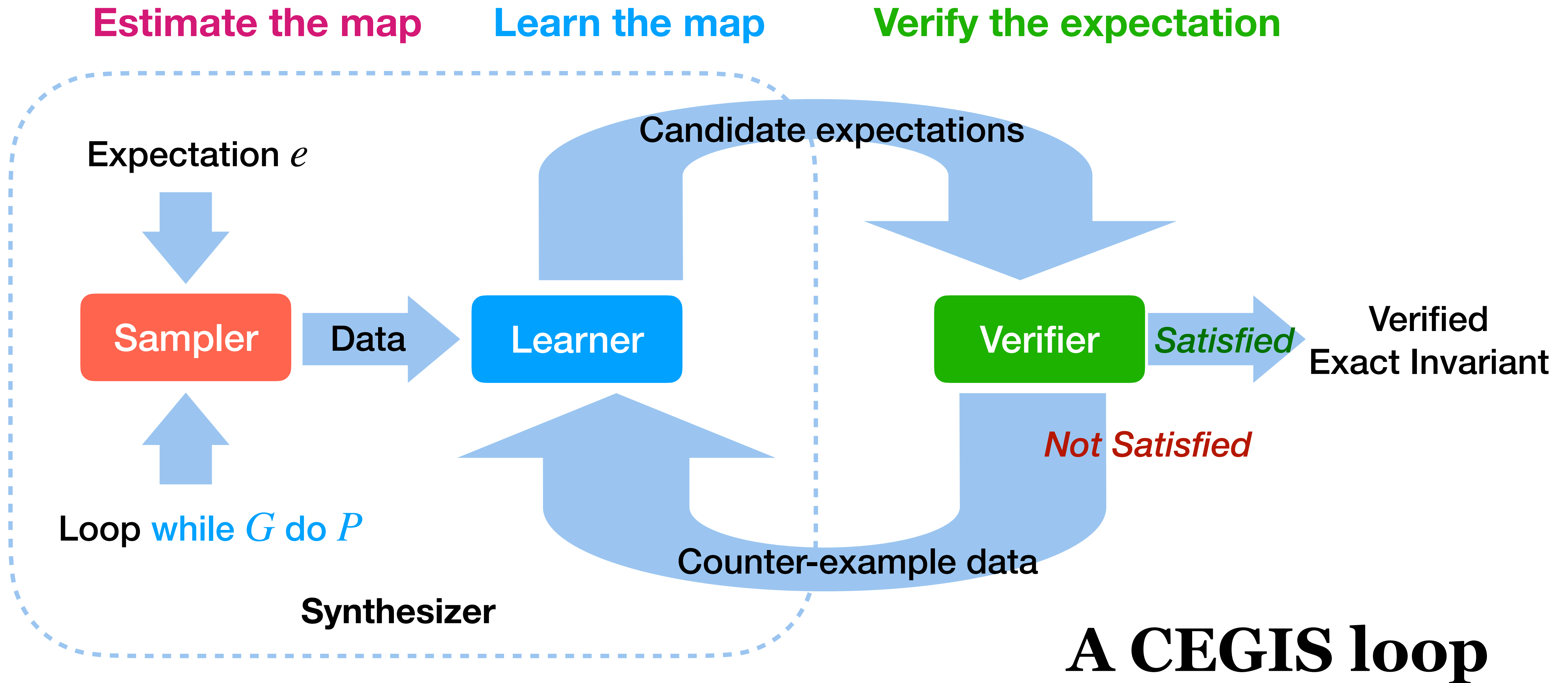
Learn the map



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Sampler

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- How to estimate $wpe(\text{while } G \text{ do } P, e)(s)$?
 - It is the **expected value** of e on the distribution obtained from running while G do P from s .
 - We can approximate expected values by **empirical means**.

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- Sample M program states in total

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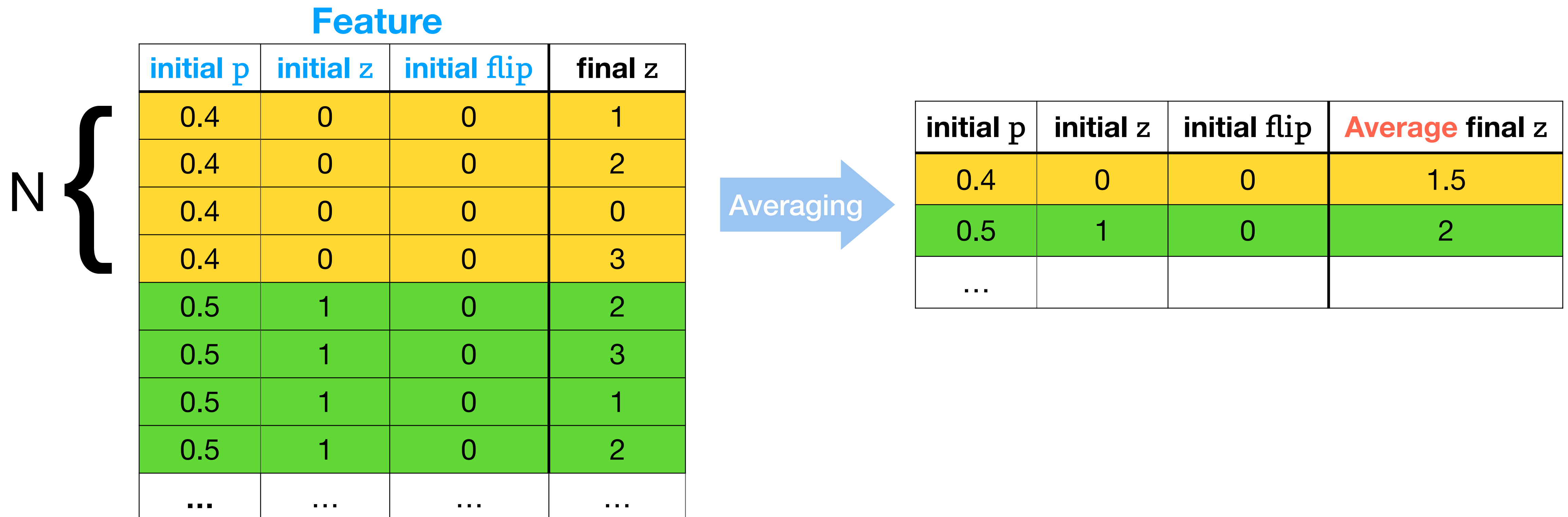
Feature

N {

initial p	initial z	initial flip	final z
0.4	0	0	1
0.4	0	0	2
0.4	0	0	0
0.4	0	0	3
0.5	1	0	2
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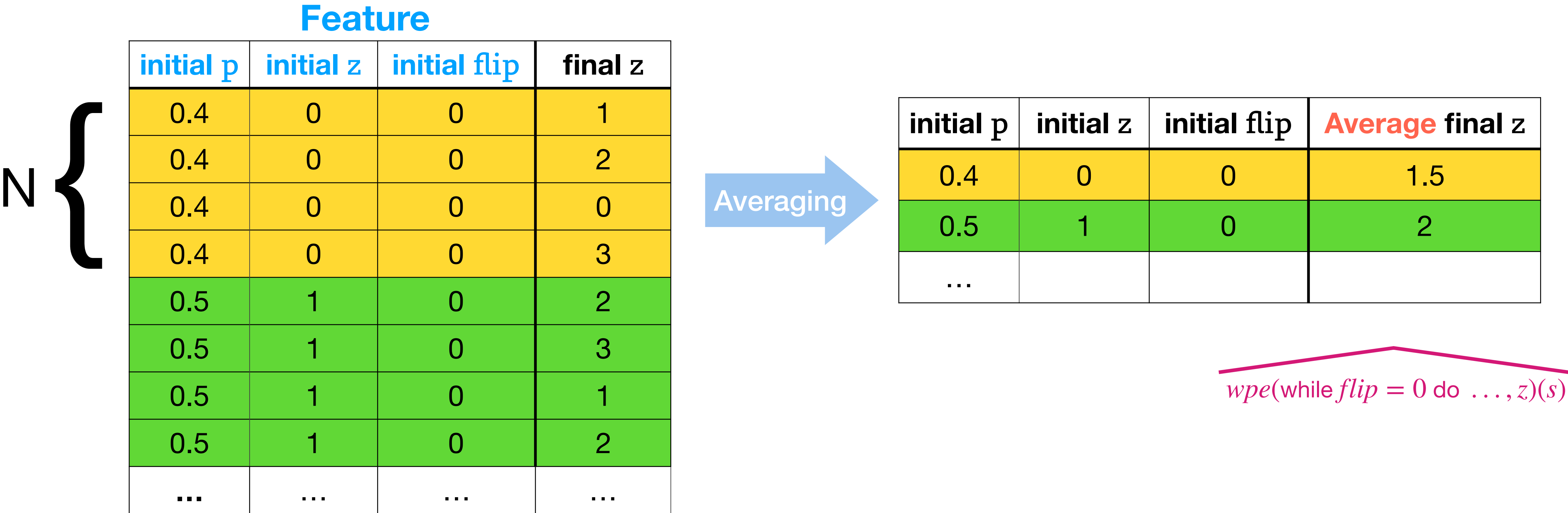
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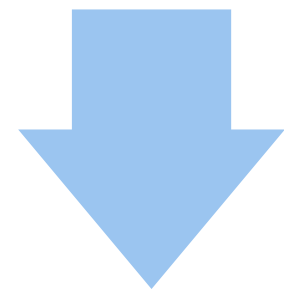
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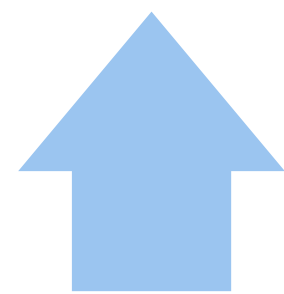


Estimate the map

Expectation e



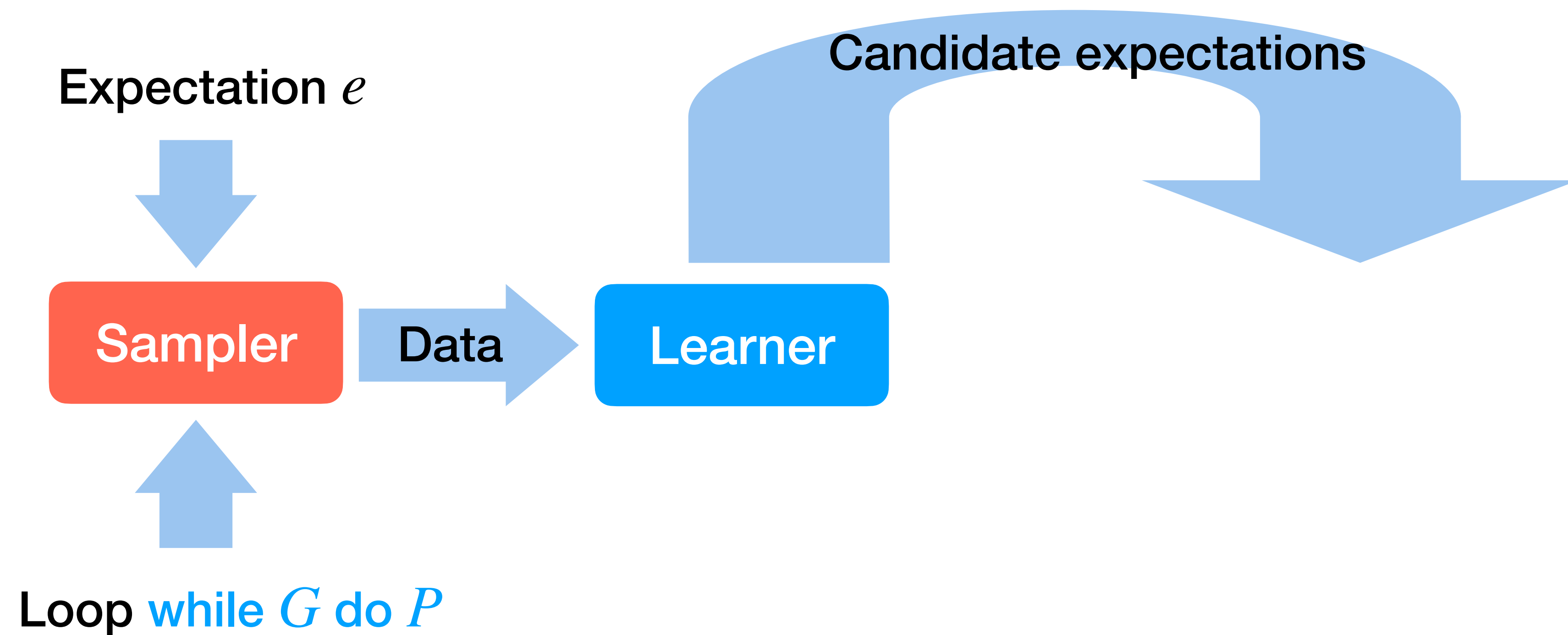
Data



Loop while G do P

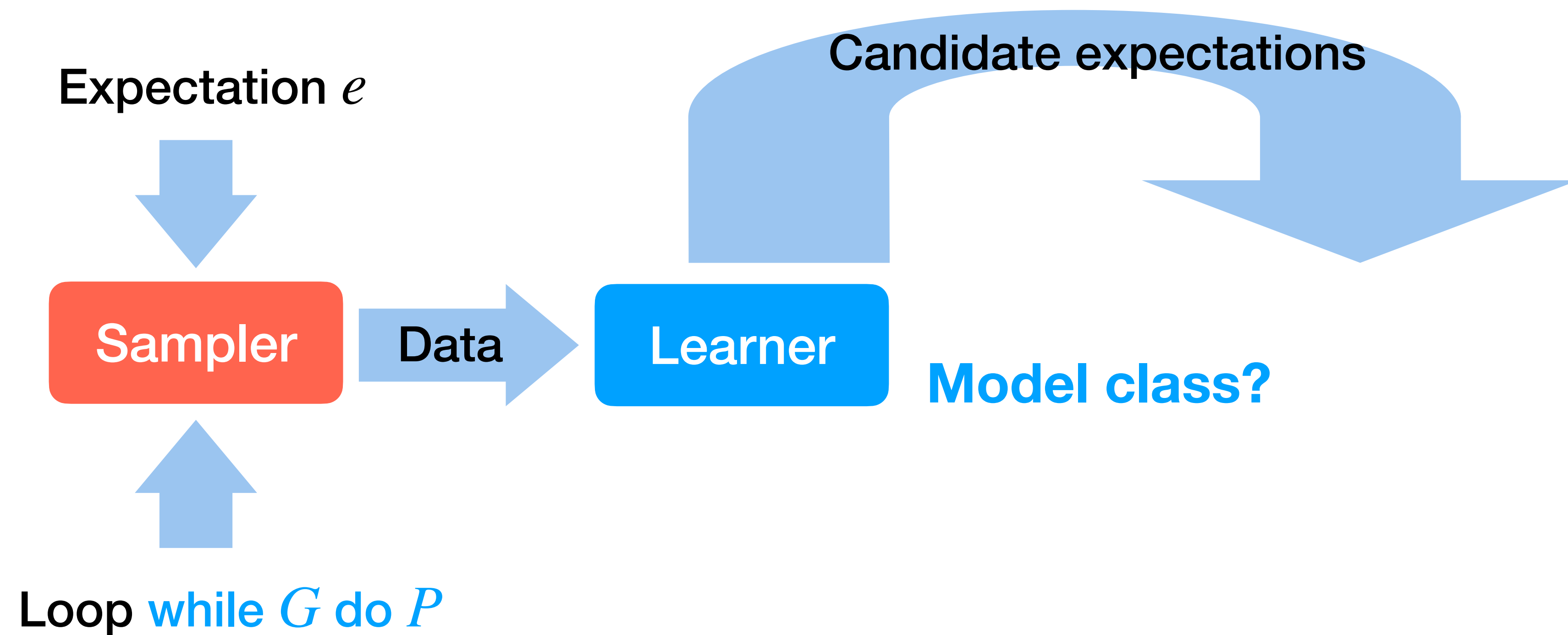
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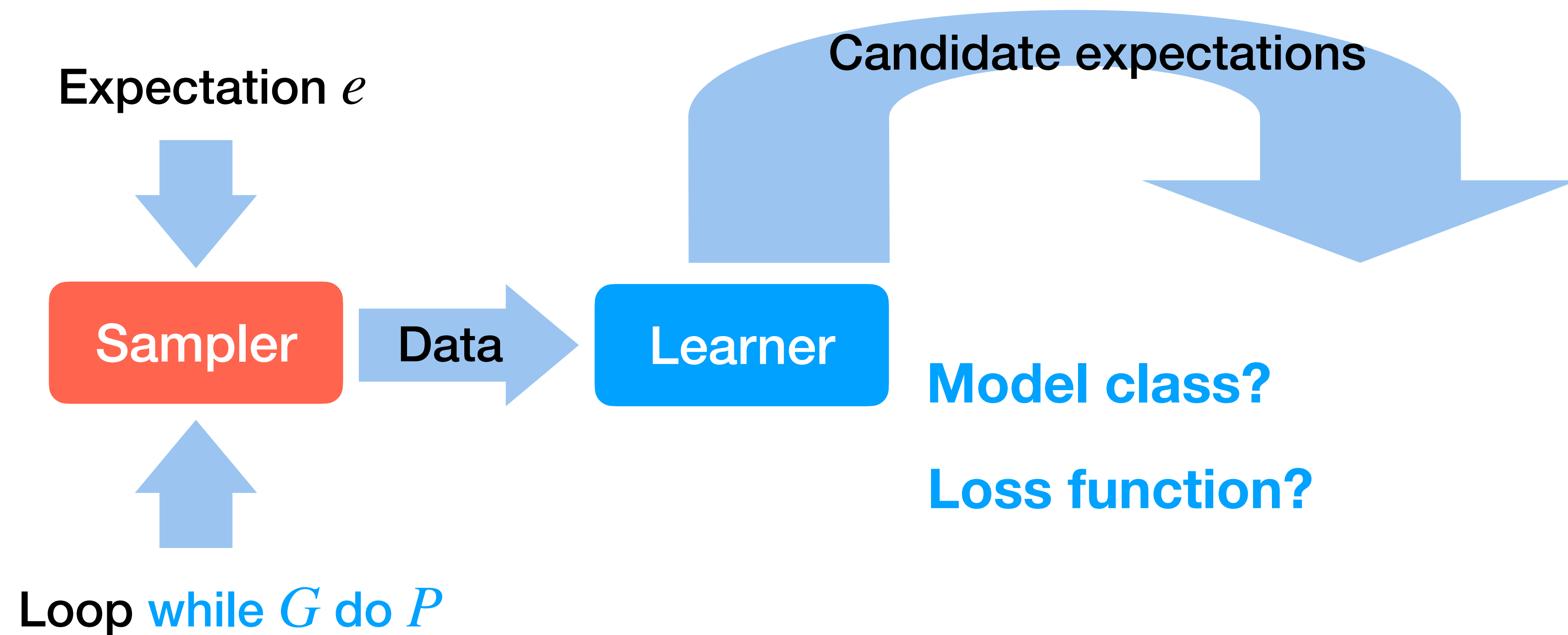
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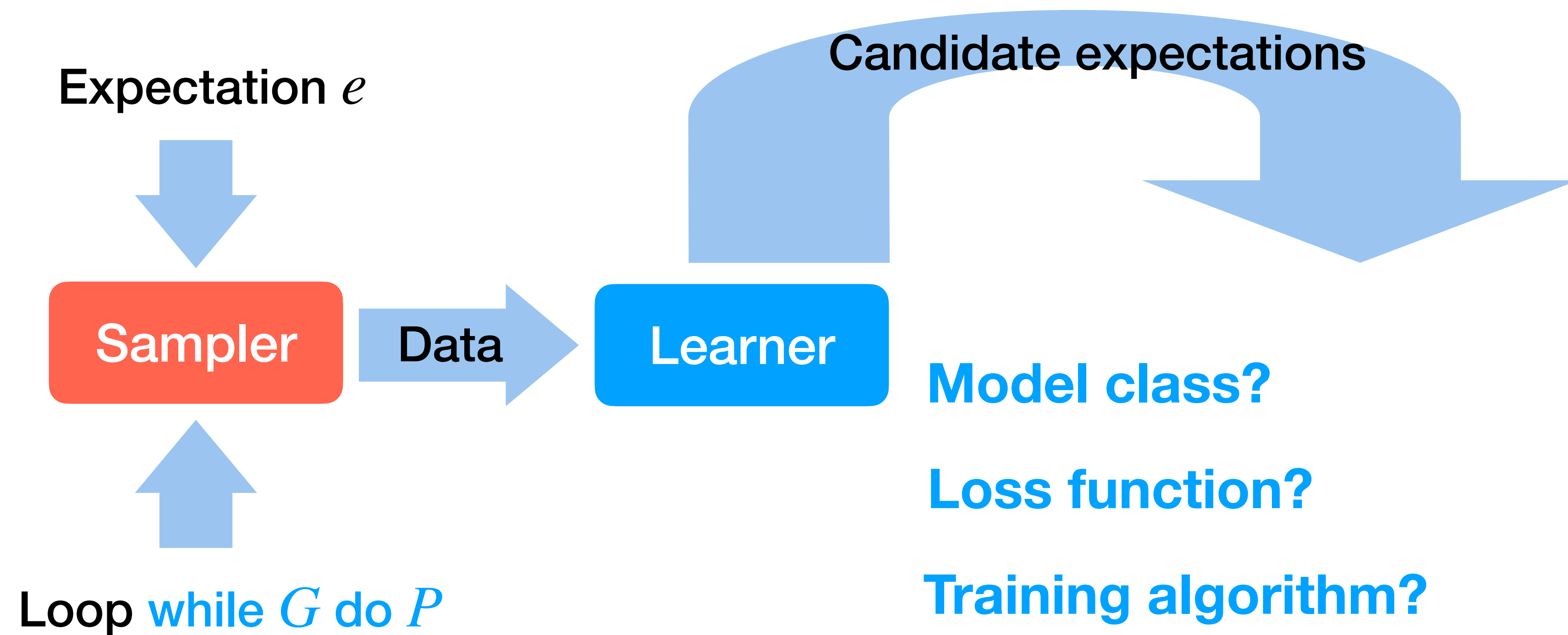
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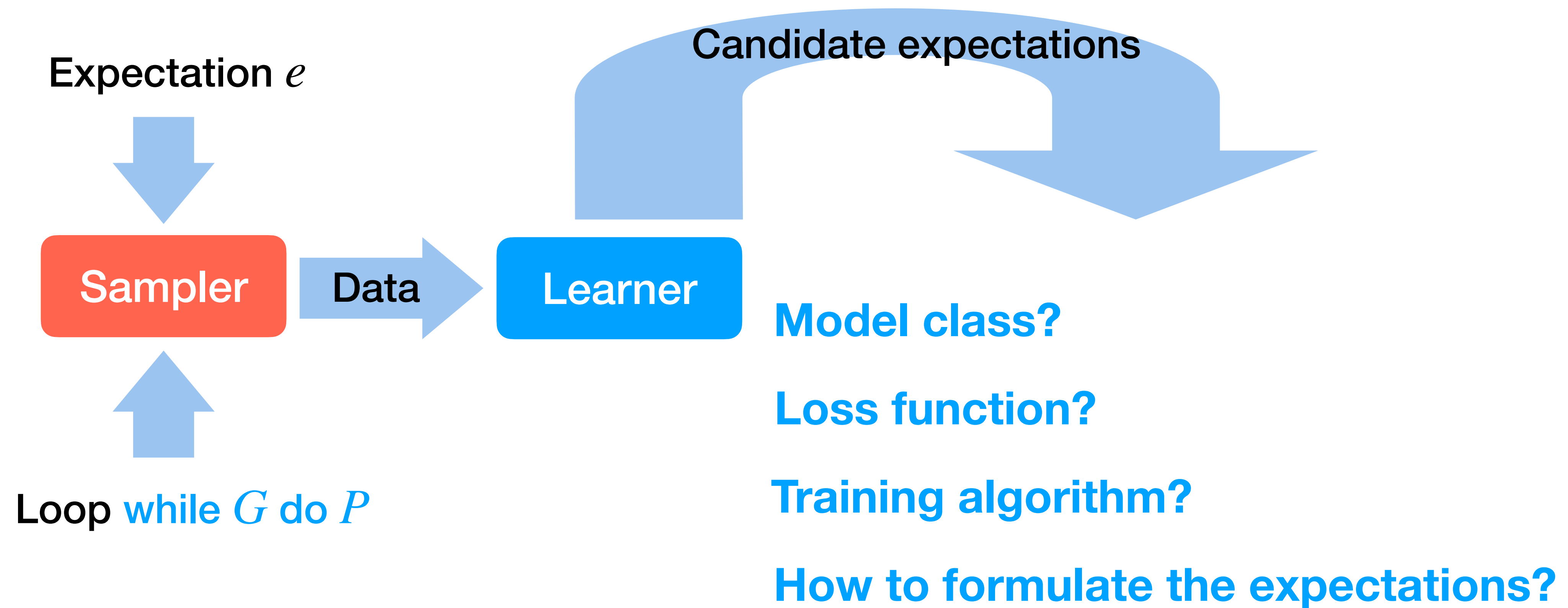
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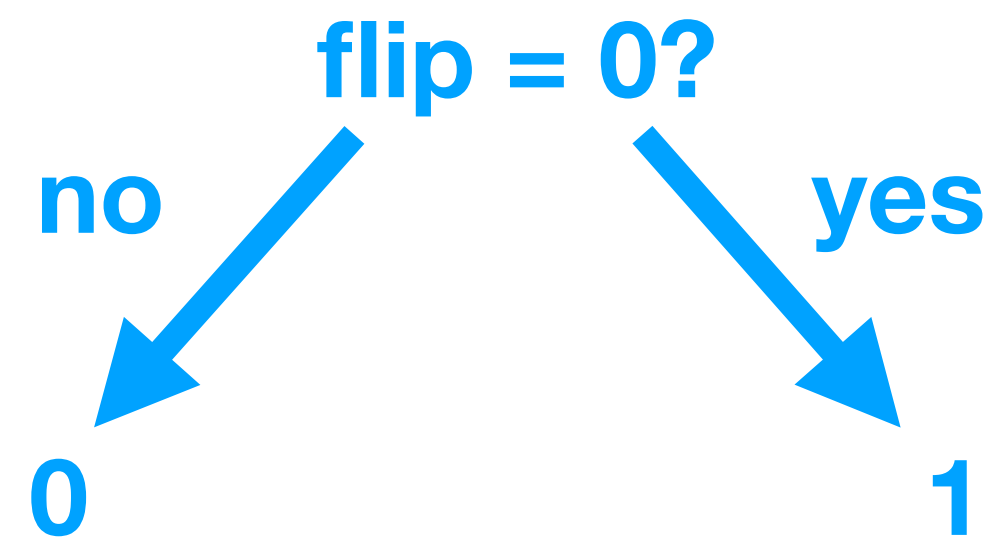
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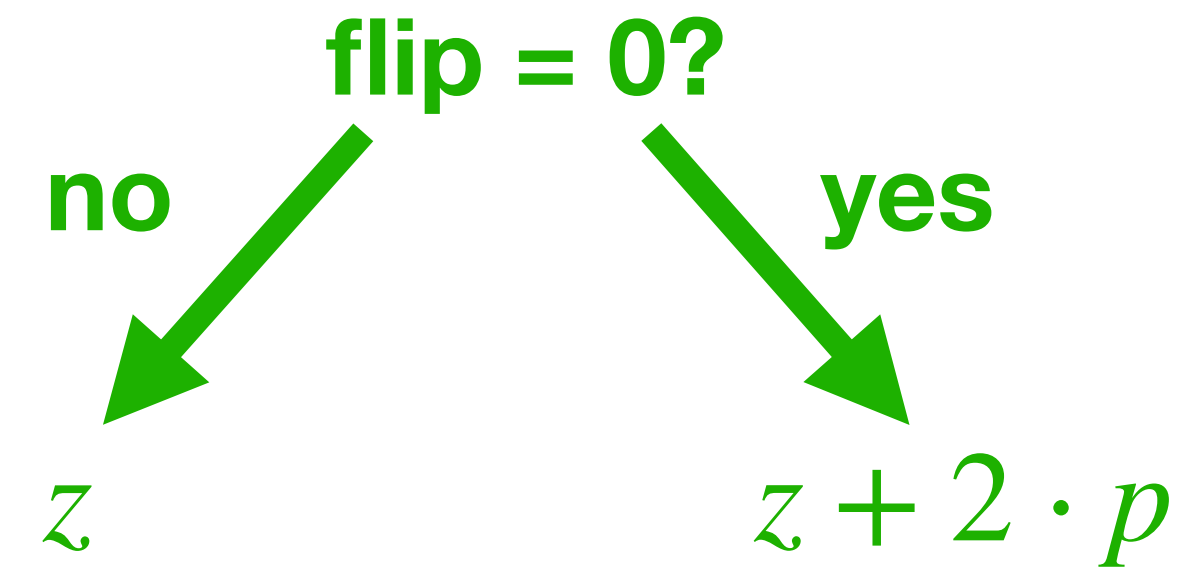
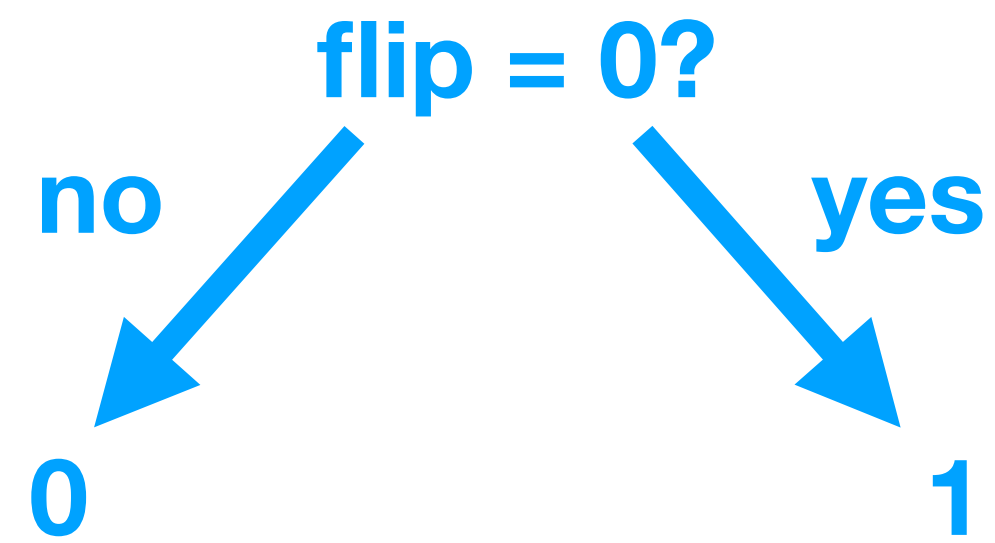
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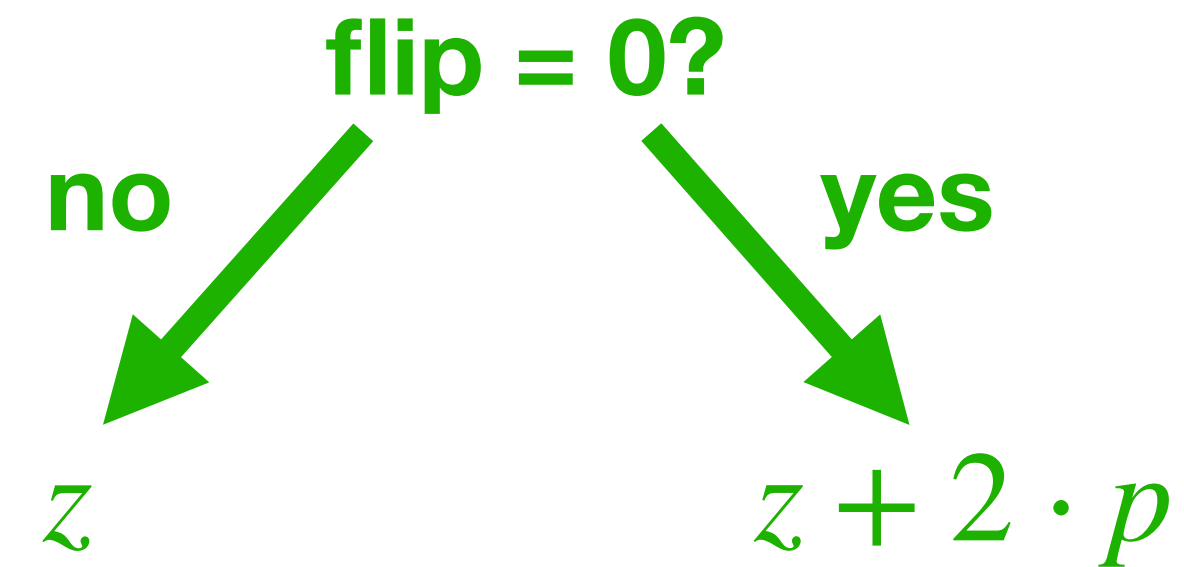
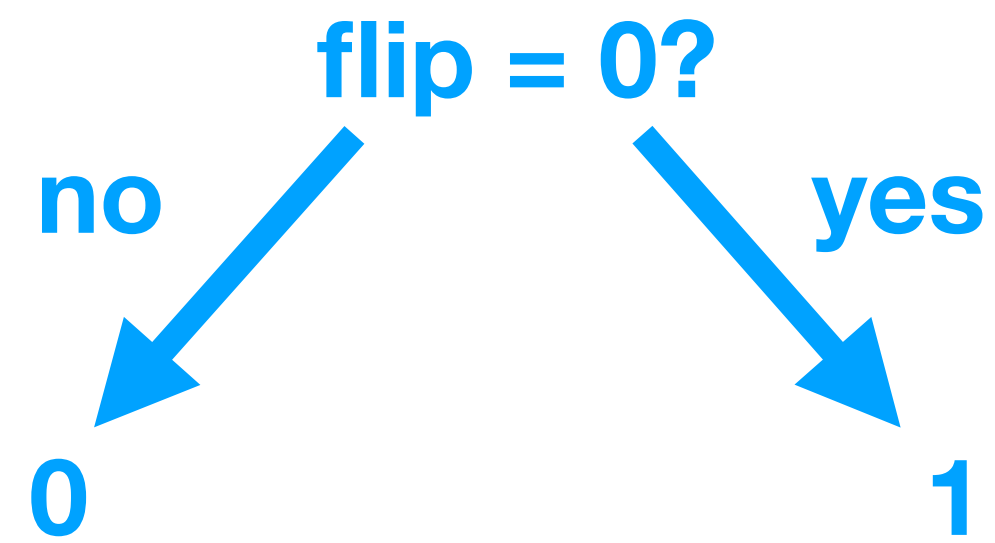
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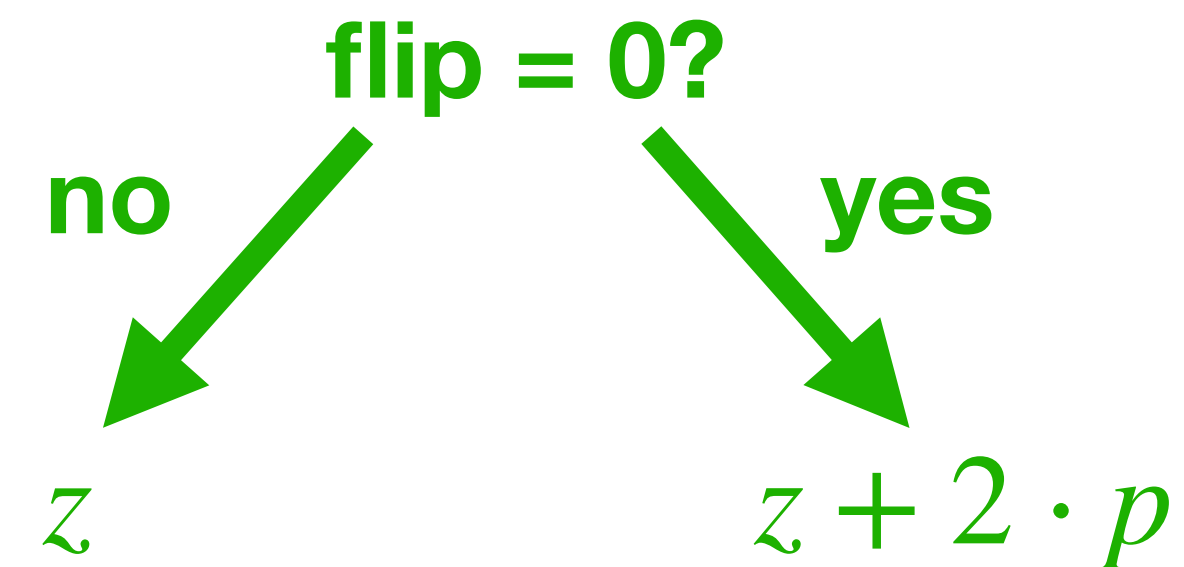
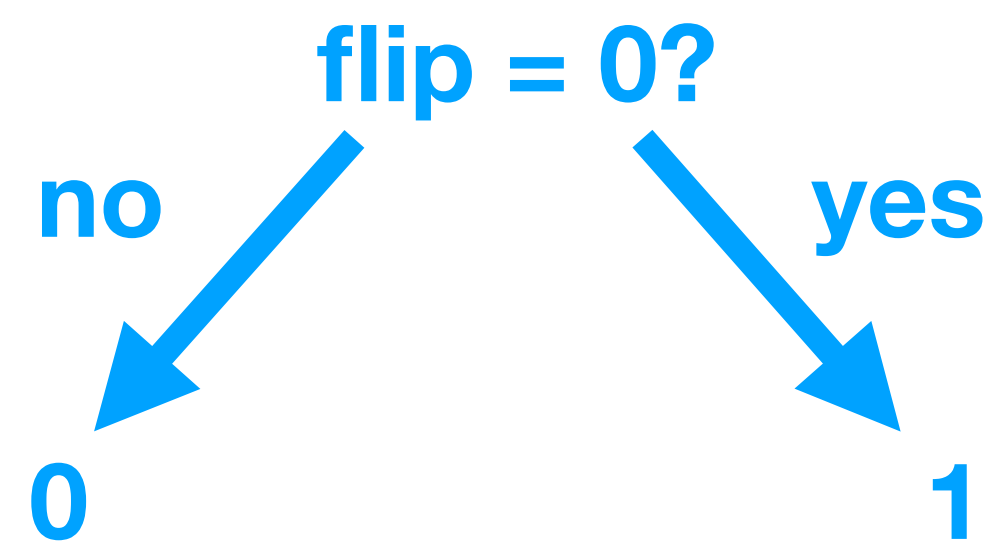


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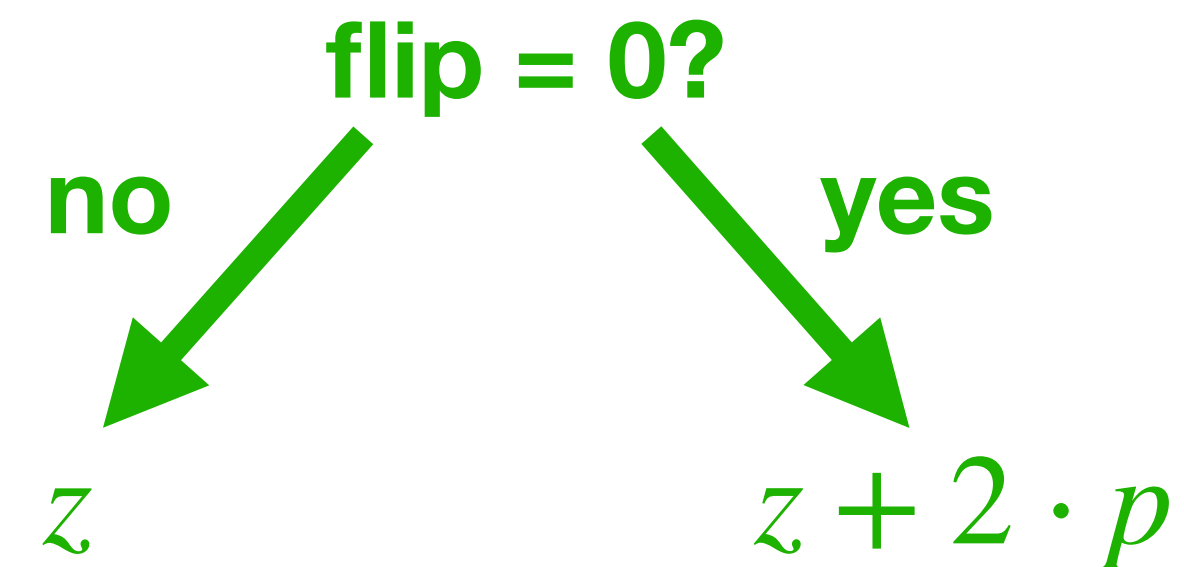
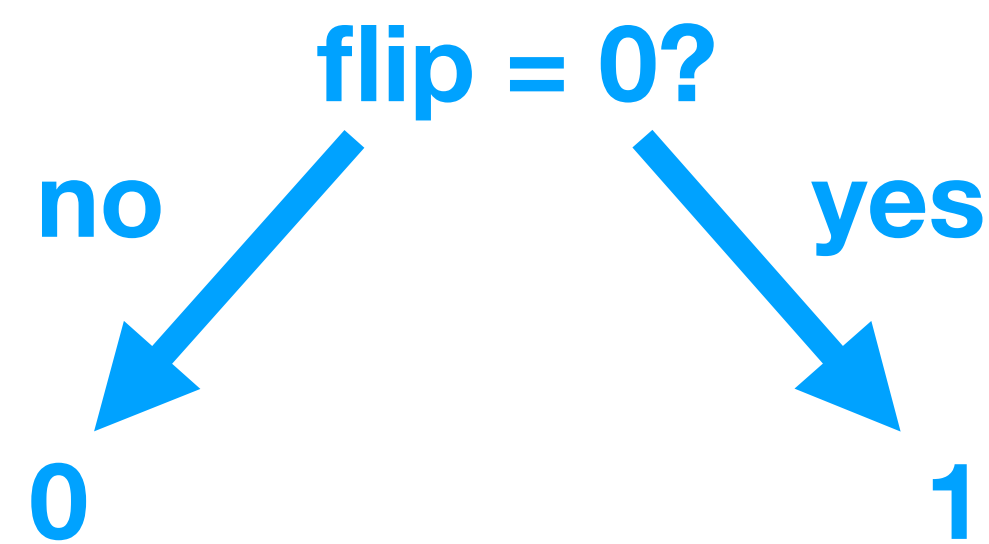


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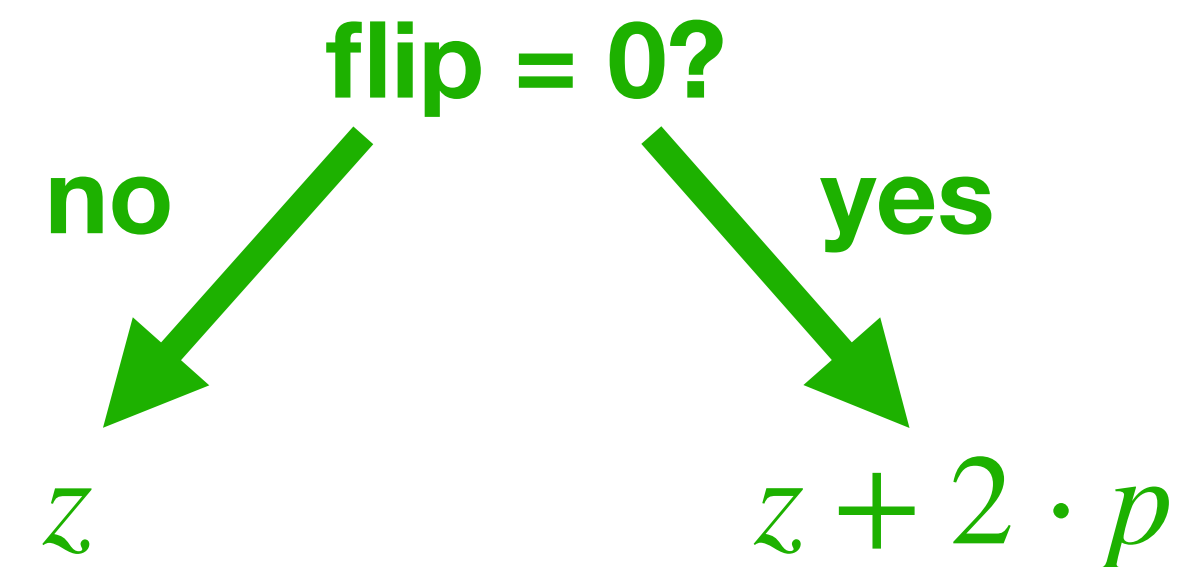
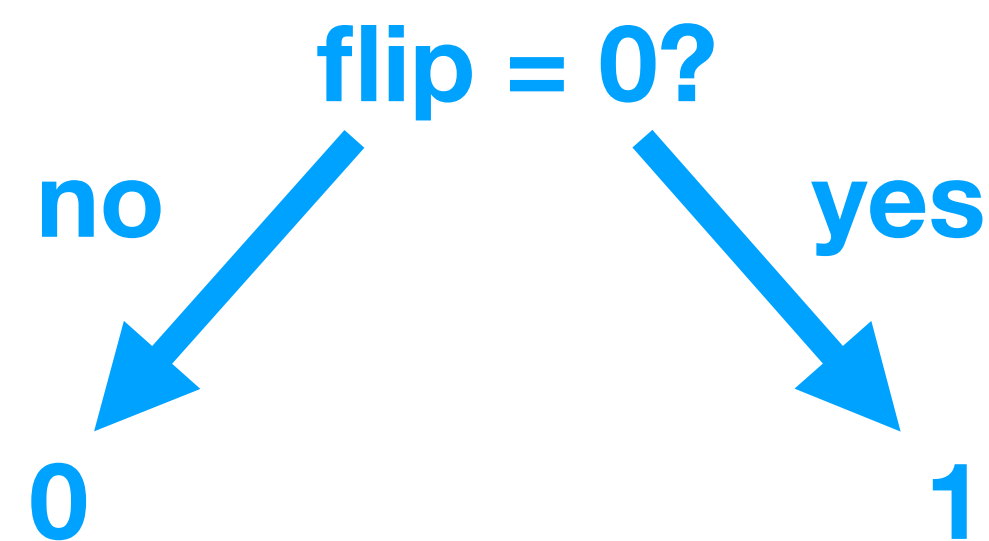


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 - Easy for verifier to manipulate

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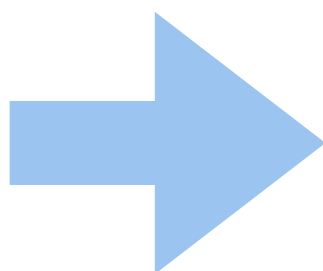
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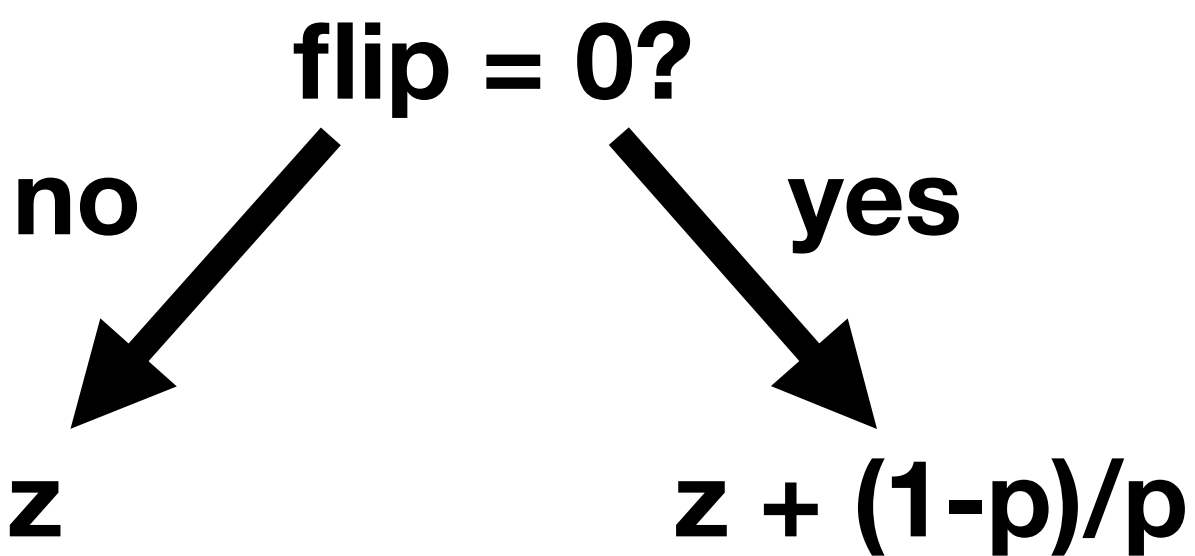
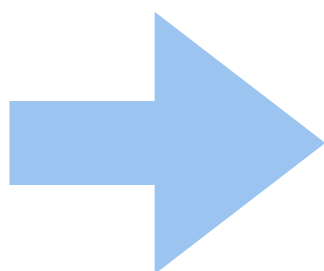
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Off-the-shelf
model tree
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Learner: Extract Expectations

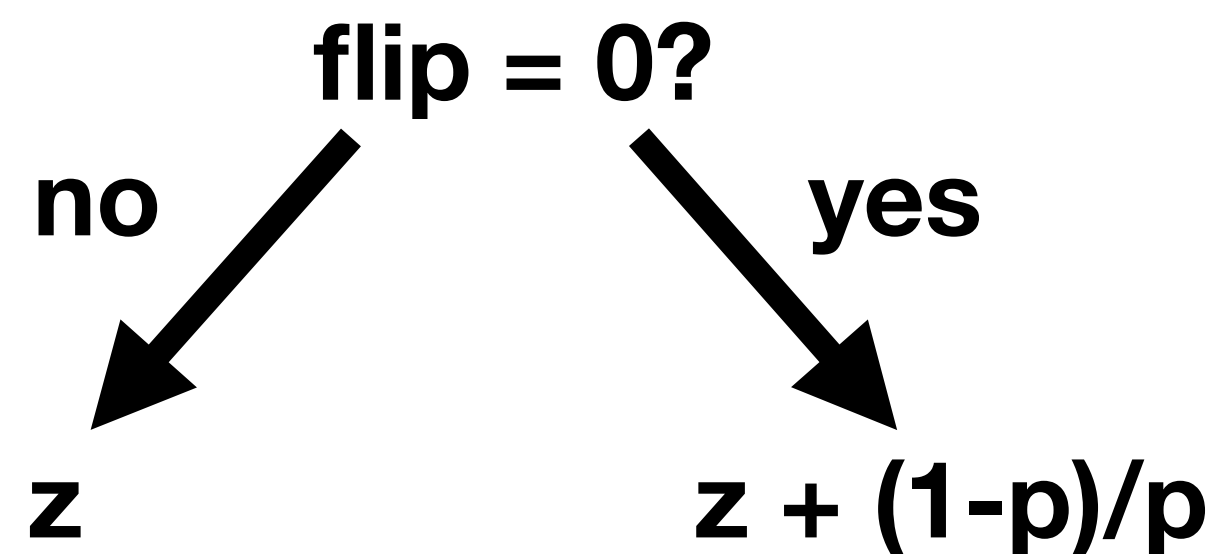
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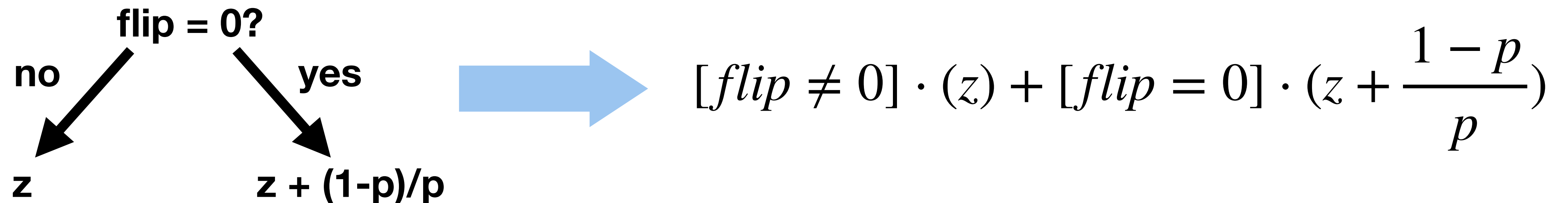
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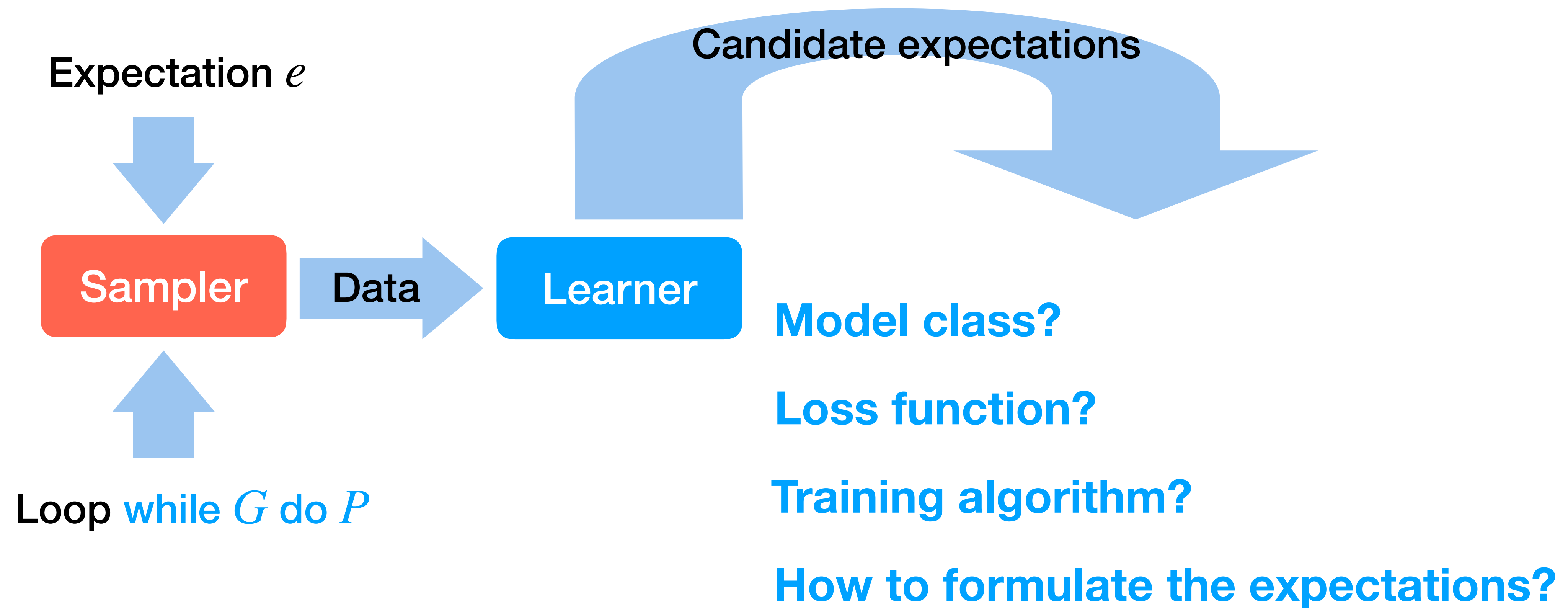
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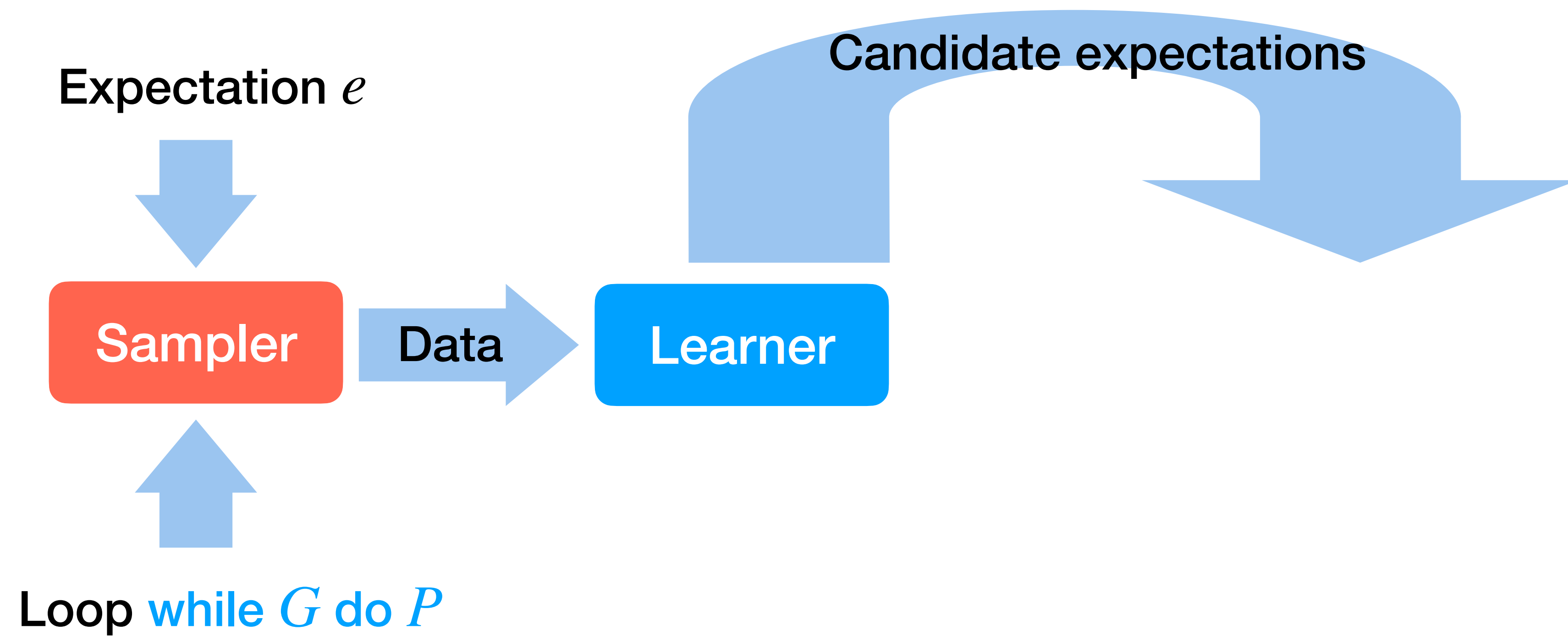
Estimate the map

Learn the map



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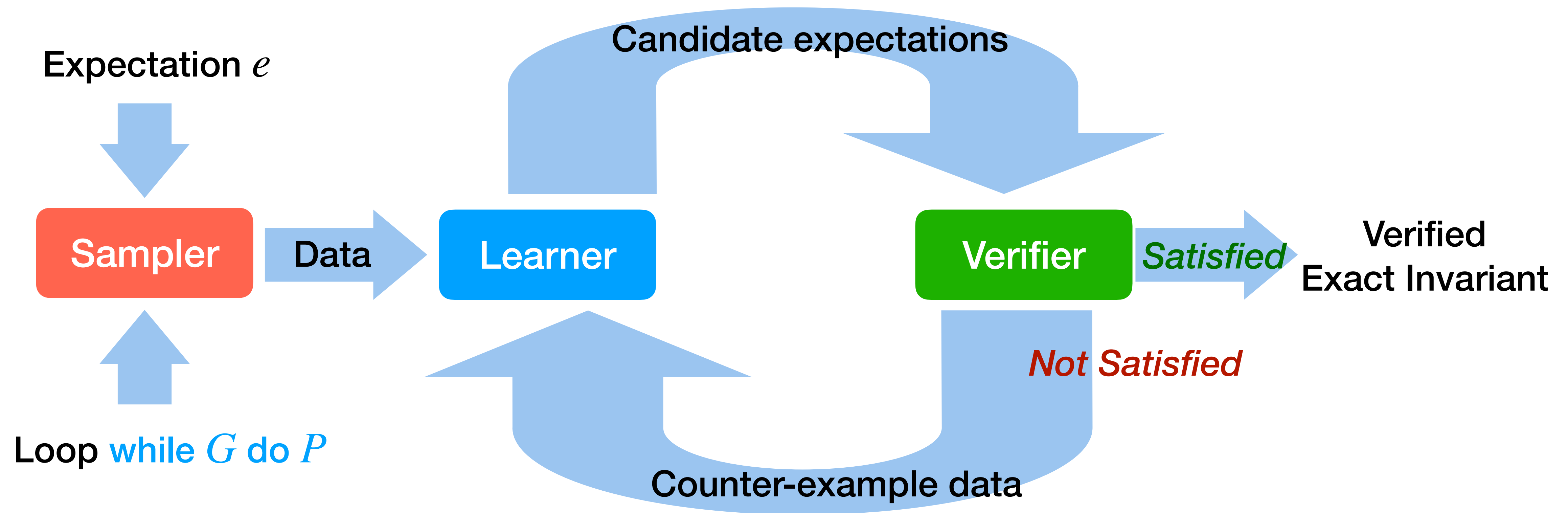
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Verifier: Check Candidate Expectations

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- Given candidate expectation I' , we use a solver to check if $I' = [G] \cdot wpe(P, I') + [\neg G] \cdot e$



Verifier: Find Counter-examples

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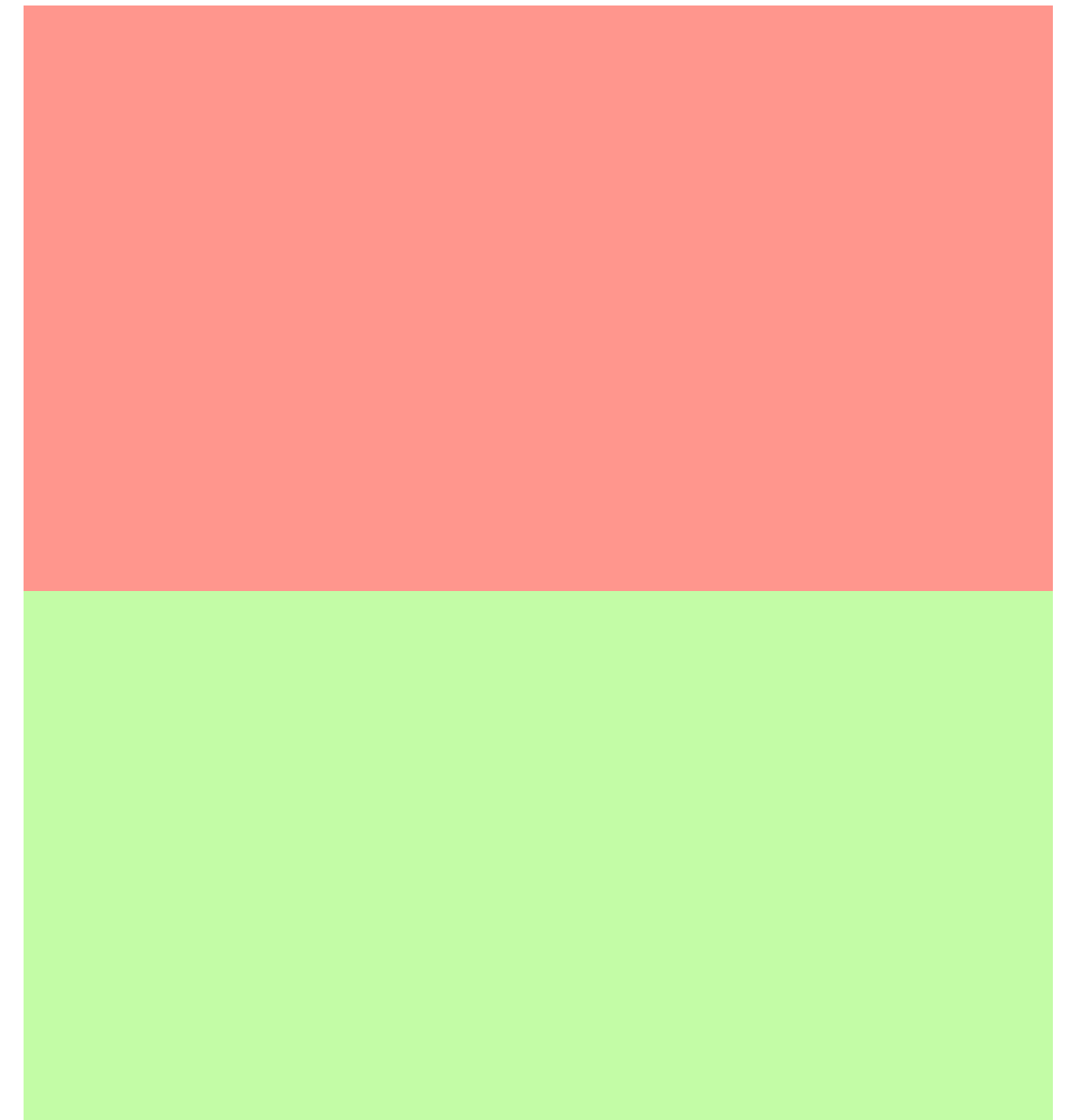
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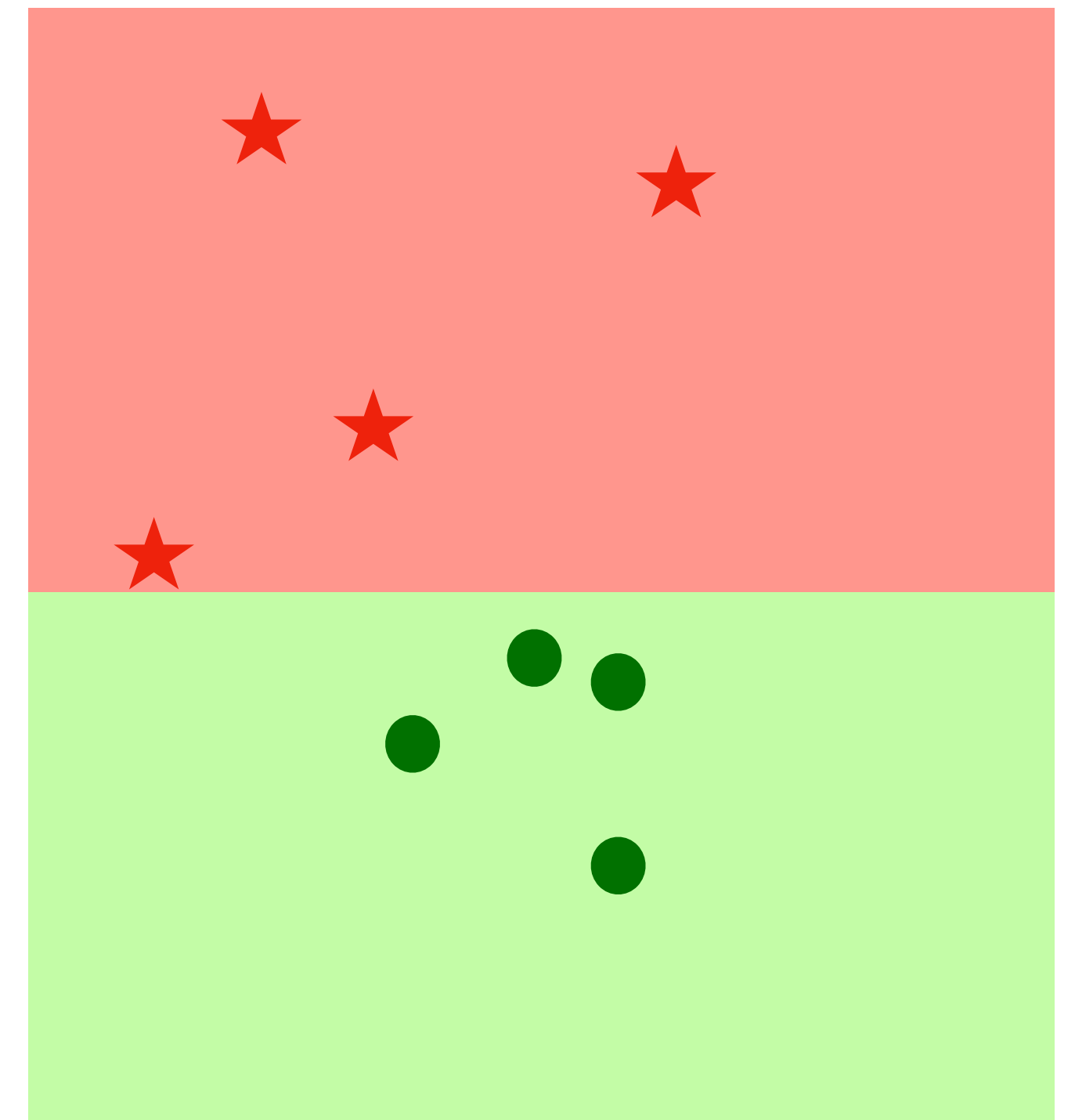
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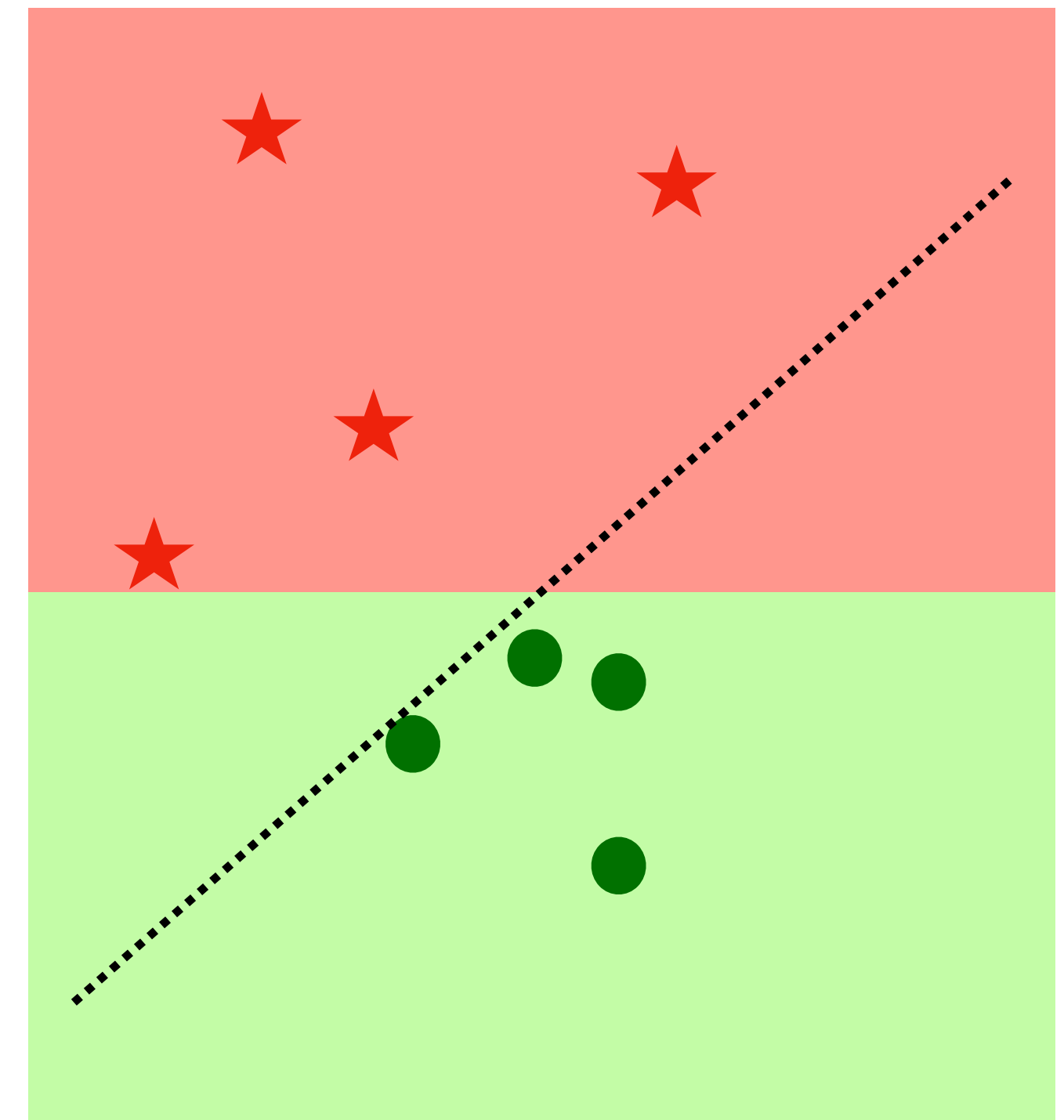
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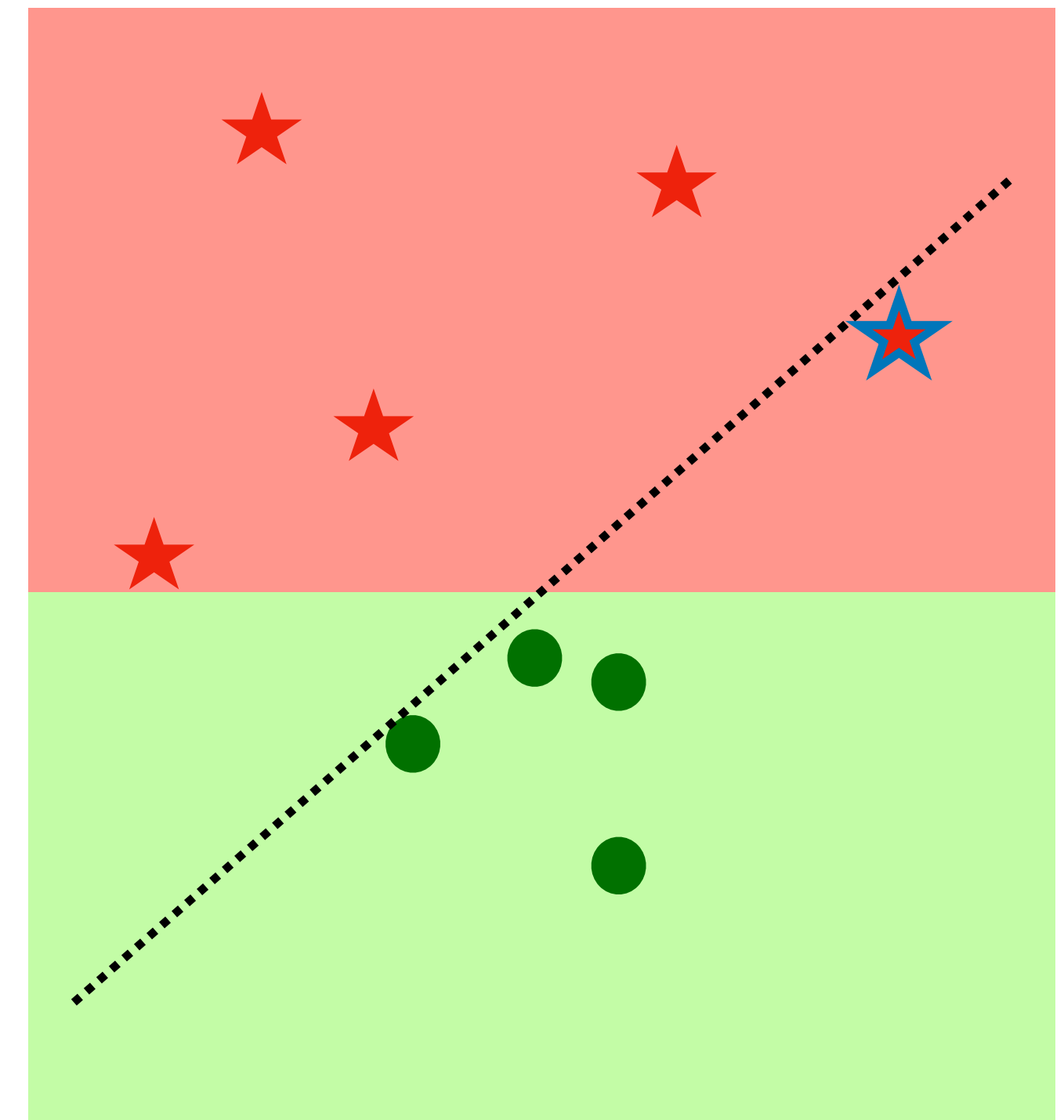
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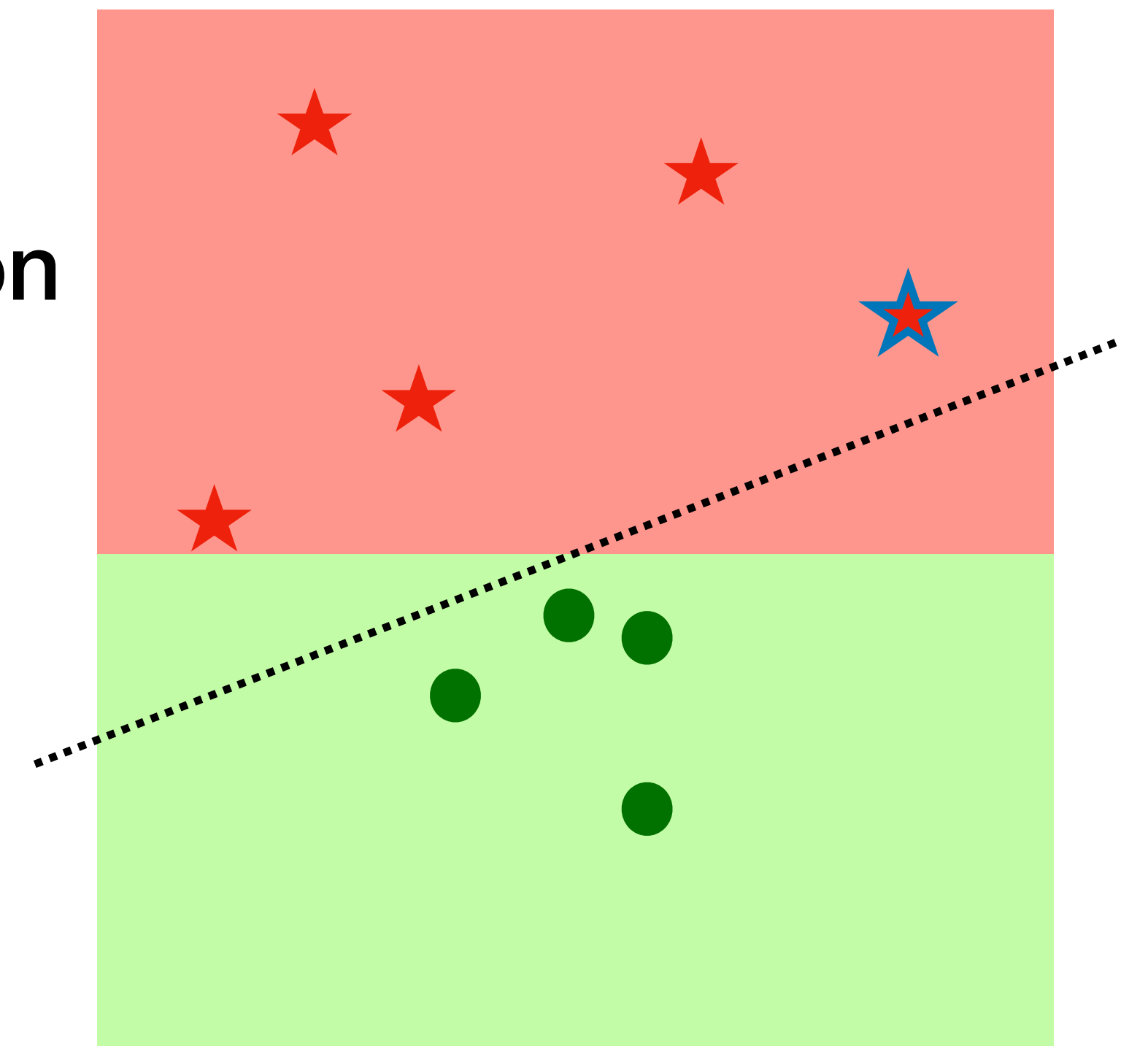
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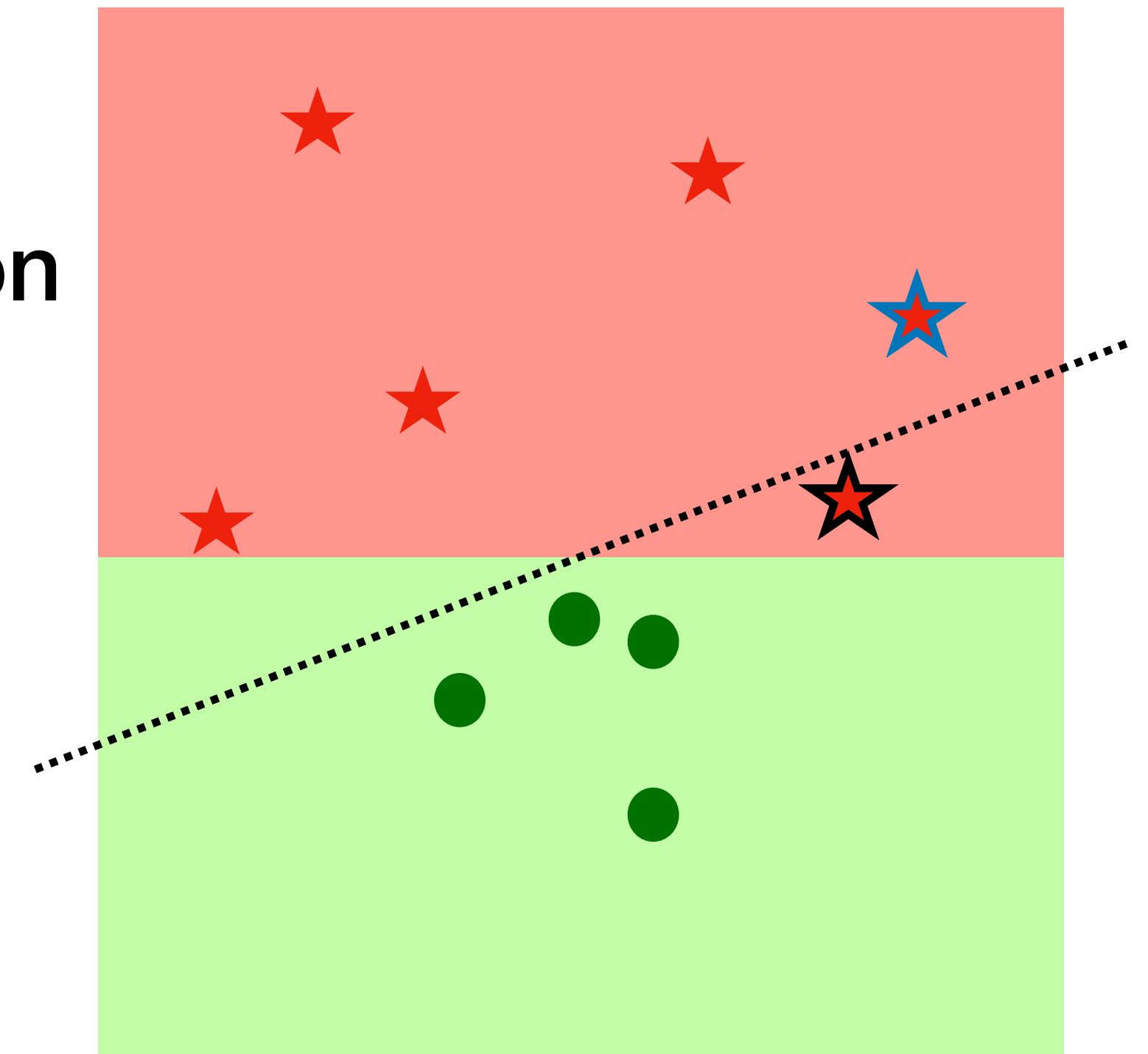
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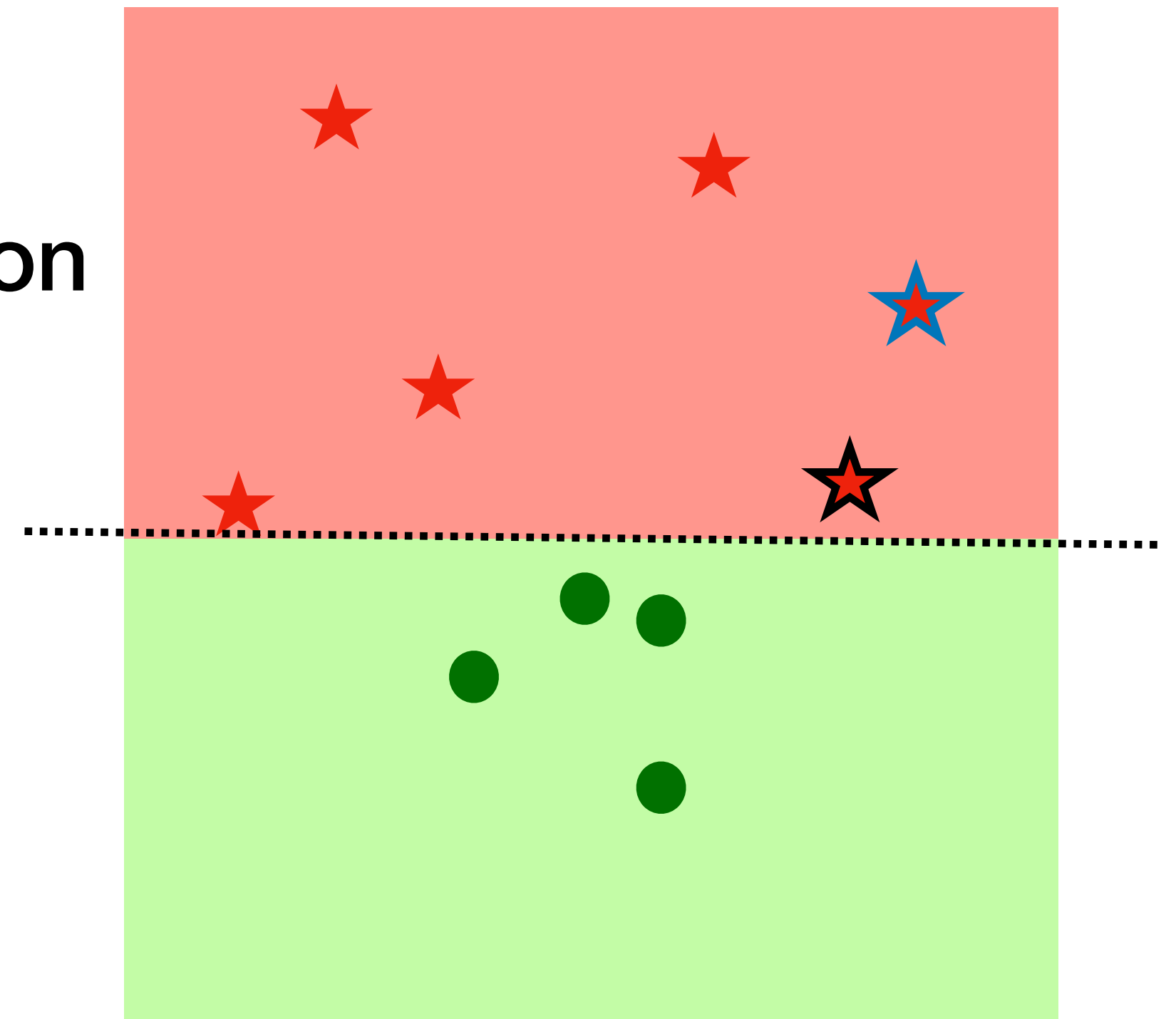
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- Too many correlated terms, e.g.,

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Not only need more data but also more powerful verifier

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- Previously, we generate **exact invariant** by learning I that approximates $wpe(\text{while } G \text{ do } P, e)$ and then check if $I = [G] \cdot wpe(P, I) + [\neg G] \cdot e$.
- Question: Can we generate **subinvariants**, i.e., I such that $I \leq [G] \cdot wpe(P, I) + [\neg G] \cdot e$?
 - Sometimes the exact invariant is too complicated
 - I is a subinvariant implies $I \leq wpe(\text{while } G \text{ do } P, e)$, not the other way

New Problem: Learning subinvariants

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- **Equivalent conditions:**

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Casting into a Learning Problem

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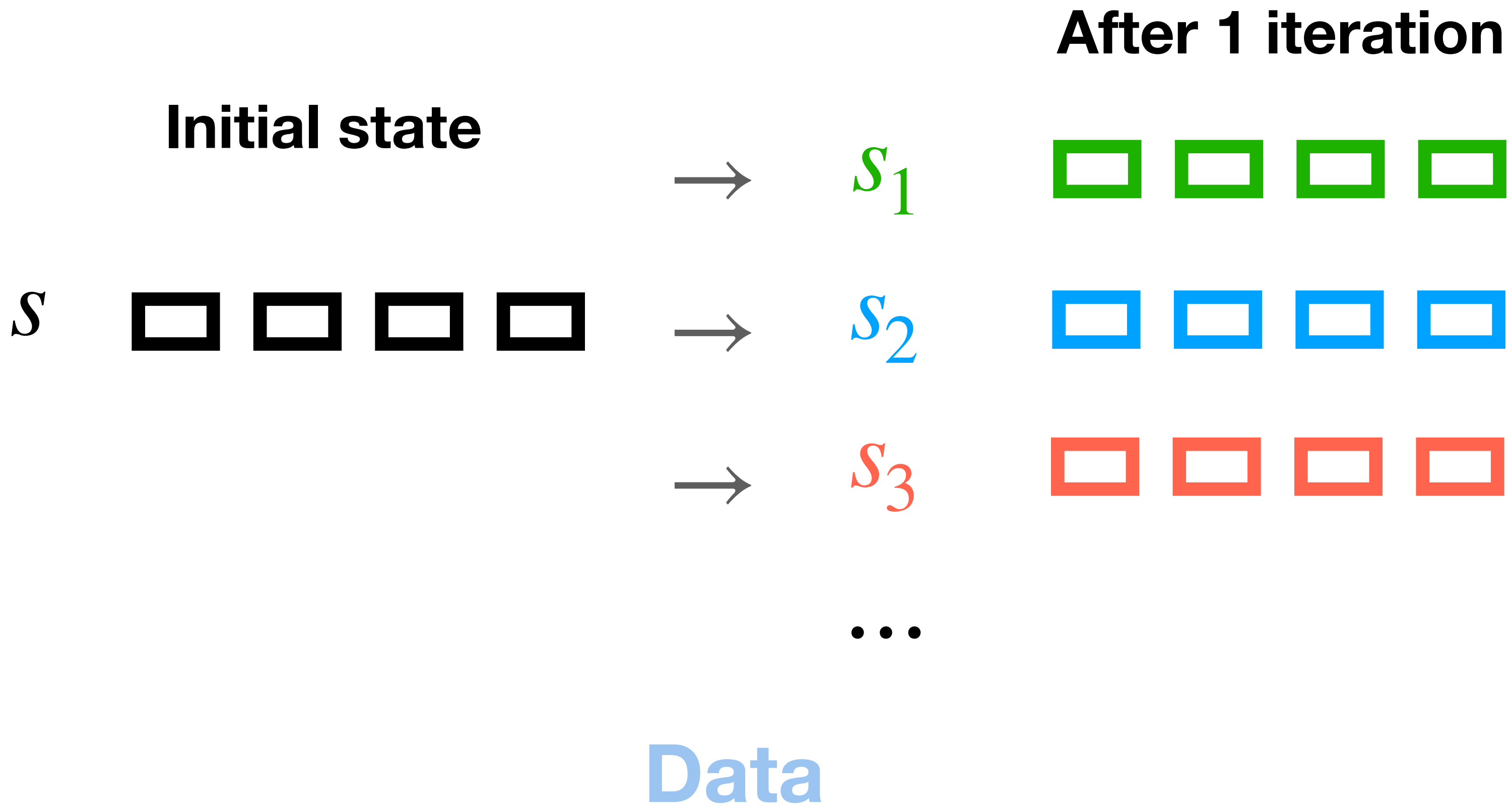
Initial state

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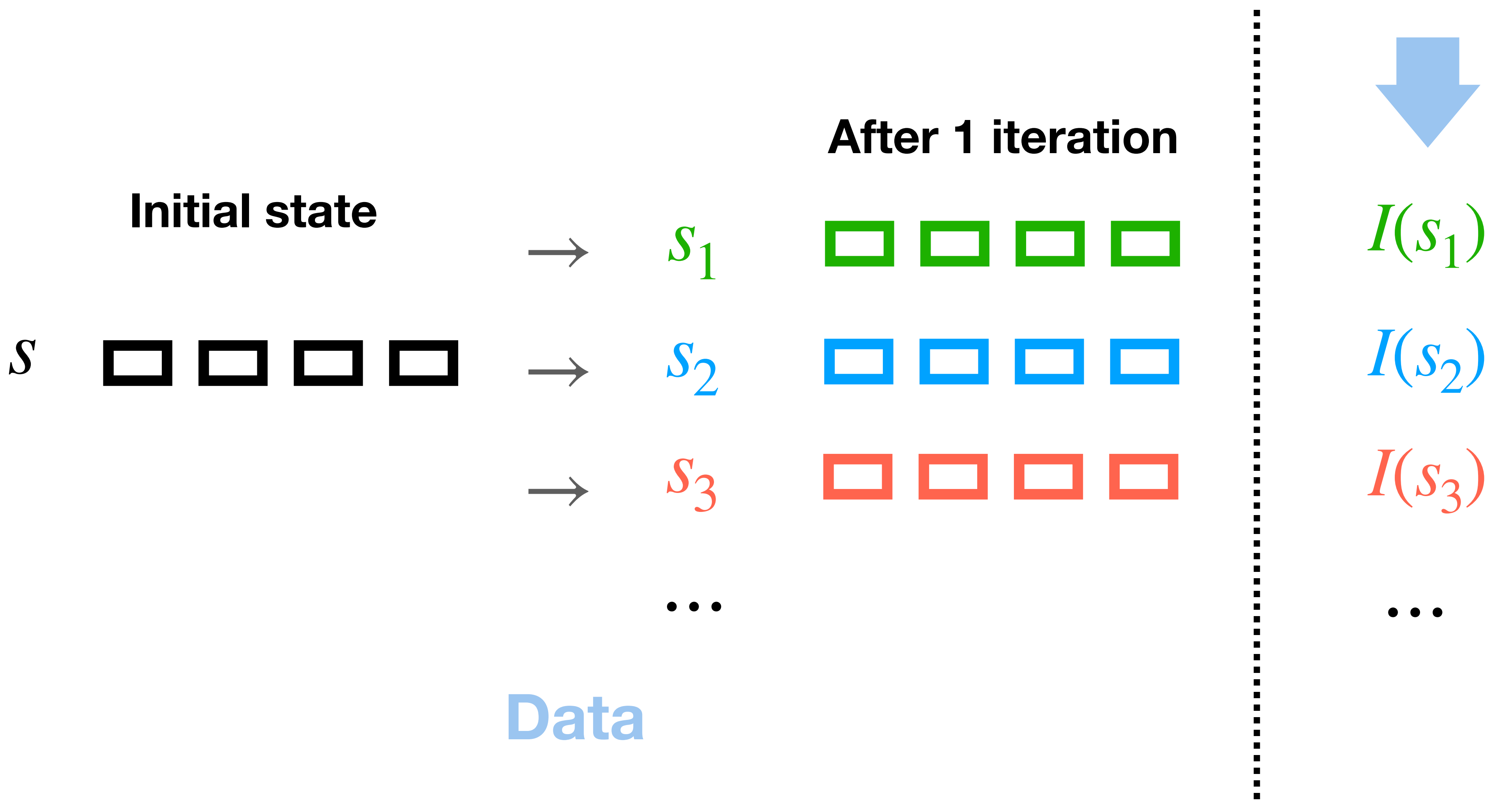


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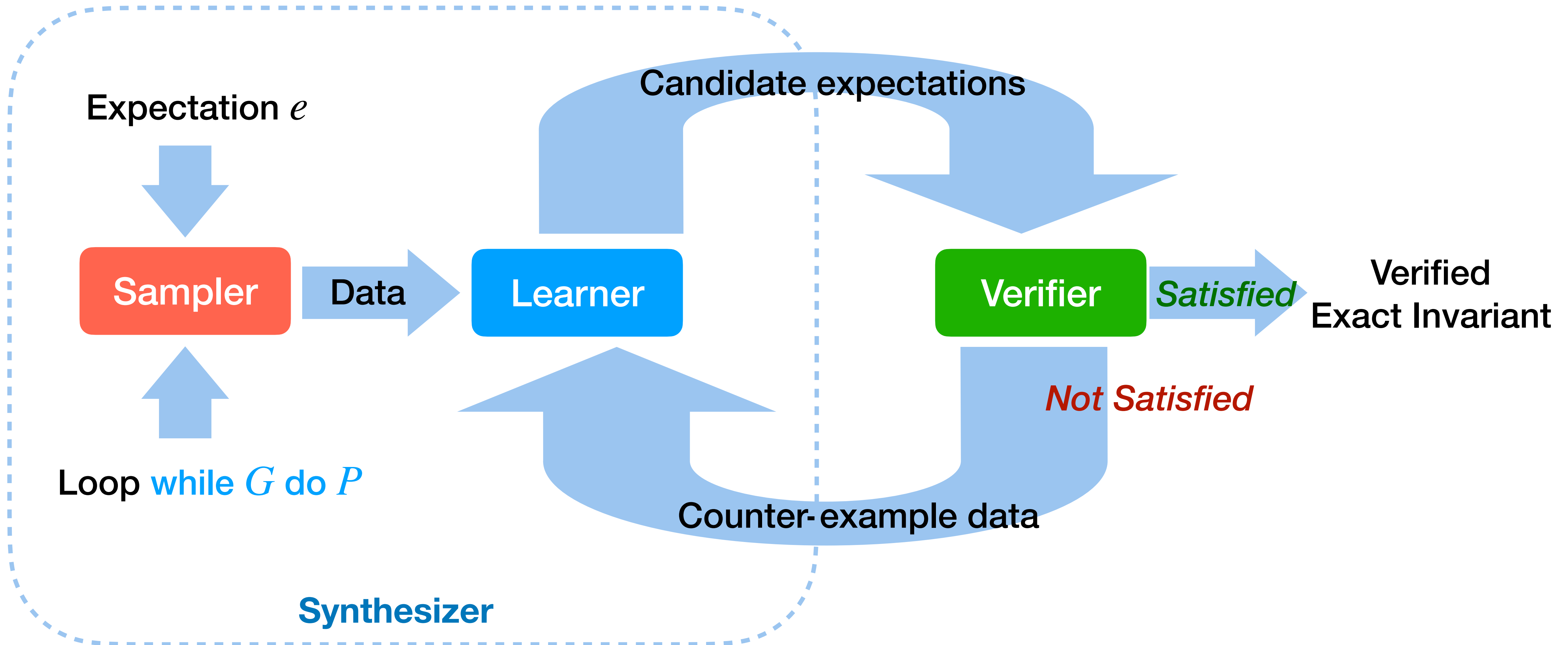
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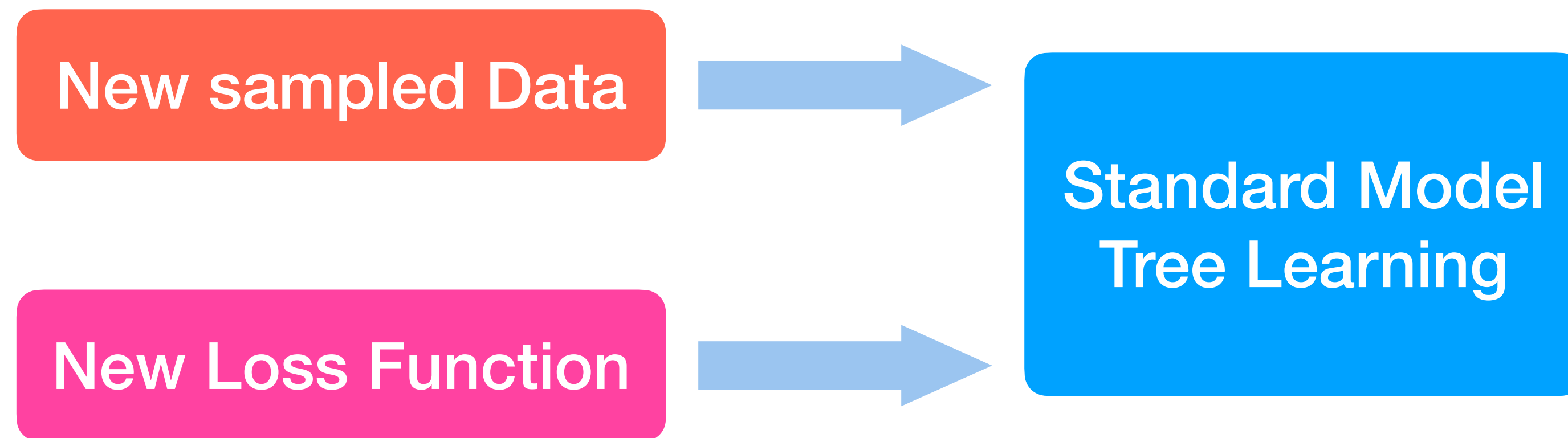


Same Method, Different Implementations

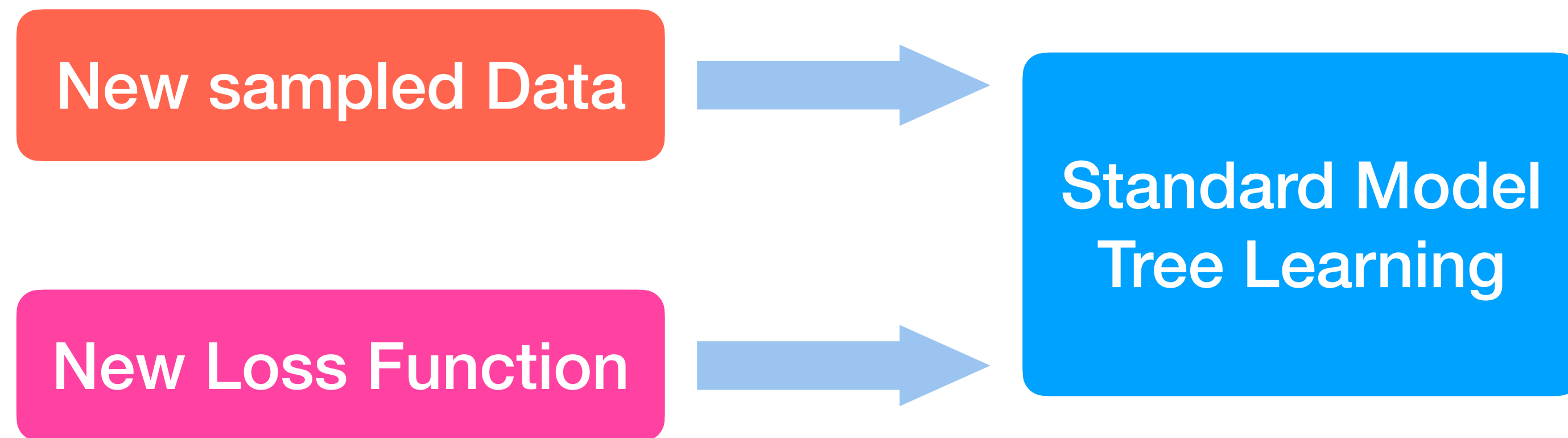


Challenge

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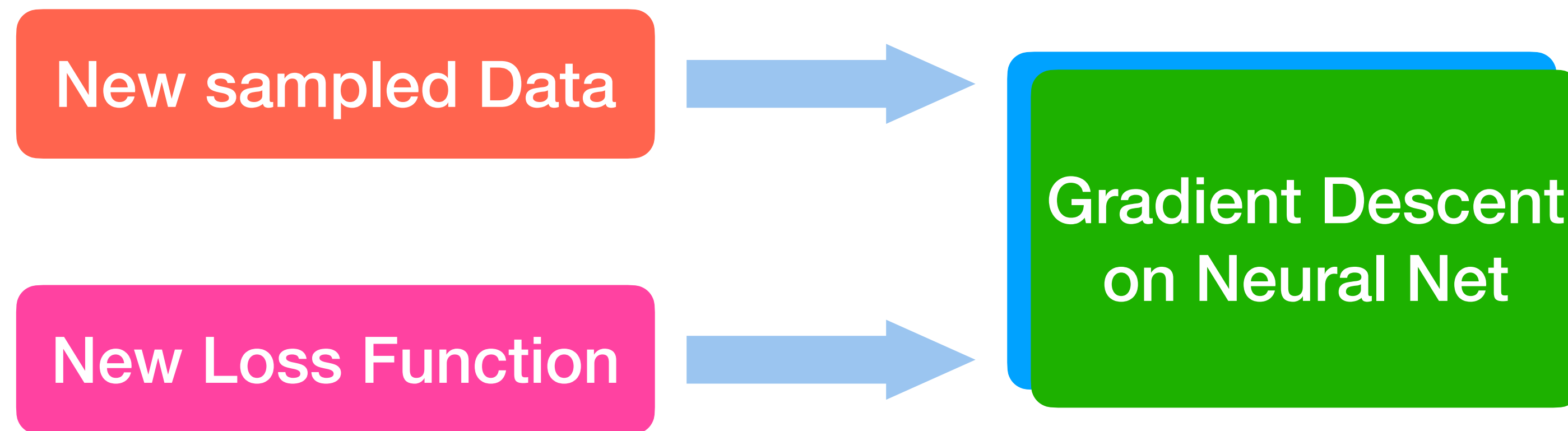


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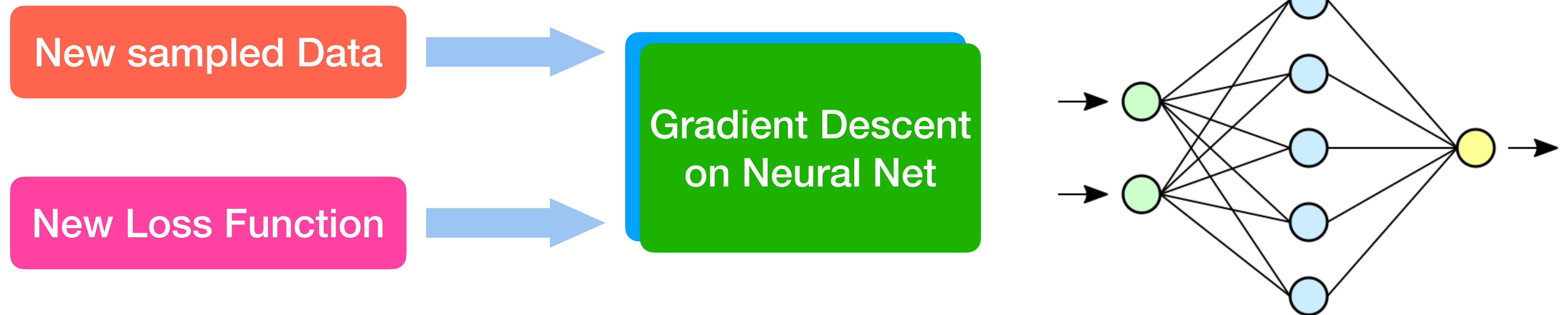


**NOT
WORKING**

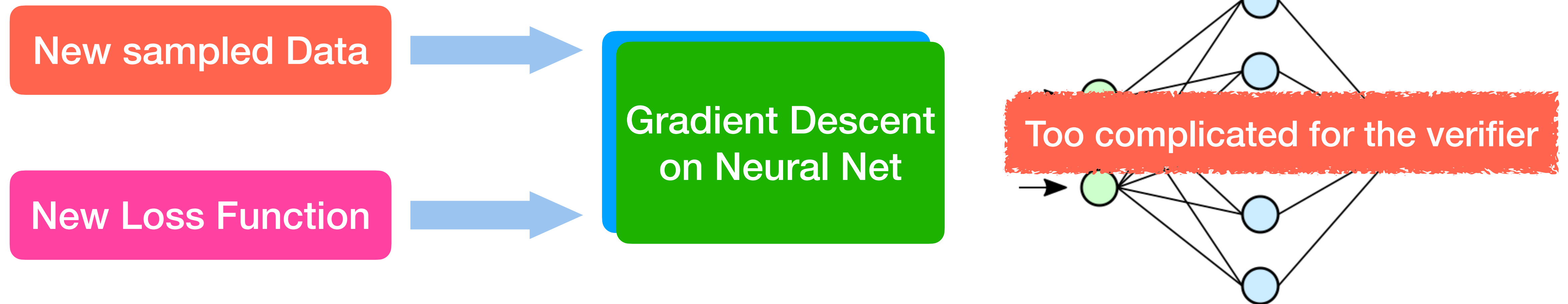
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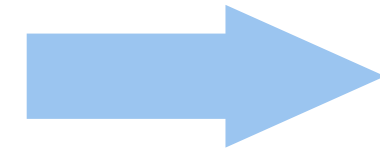


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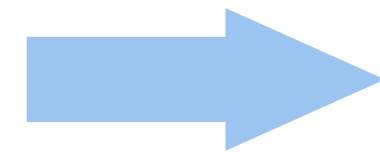


Our Solution

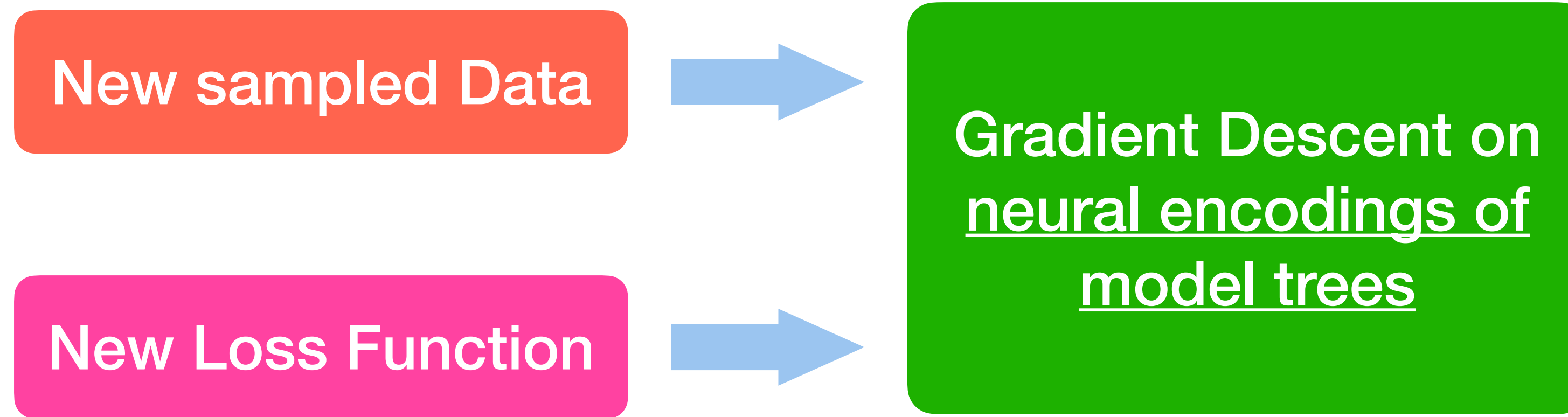
New sampled Data



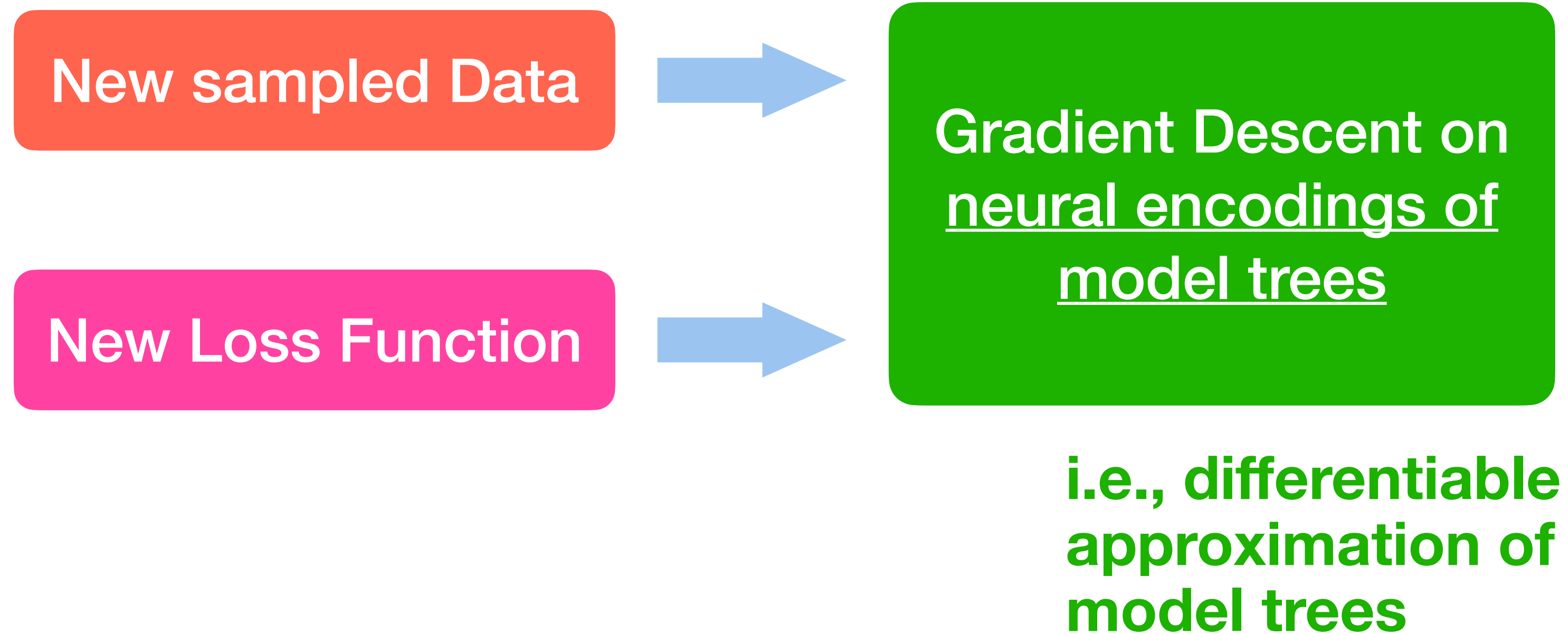
New Loss Function



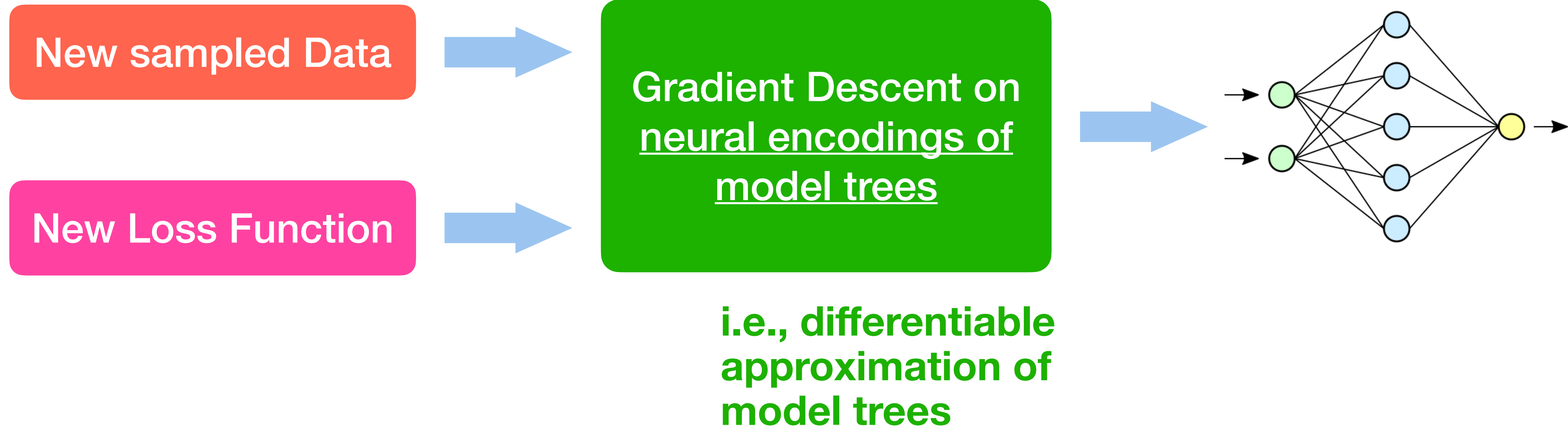
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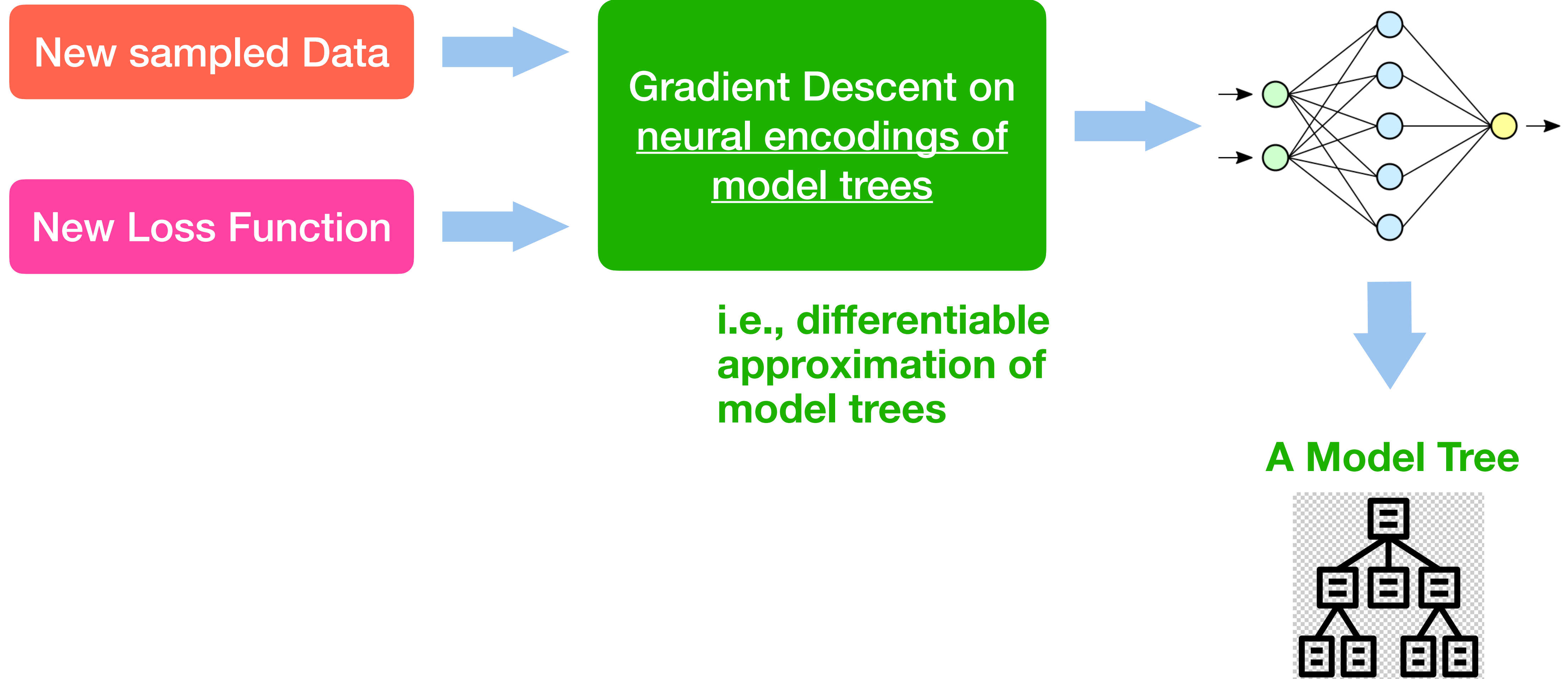
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 - It seems a hard learning program.

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- Our learning problem in exact invariant generation is almost the same.
- Our learning problem in subinvariant generation is a bit more general.

The State of Art of Symbolic Regression

Table 1: Recovery rate of several algorithms on the Nguyen benchmark problem set across 100 independent training runs. Results of our algorithm are obtained using PQT; slightly lower recovery rates were obtained using VPG and RSPG training (see Table 3 for comparisons).

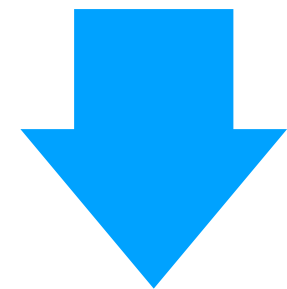
Benchmark	Expression	Recovery rate (%)					
		Ours	DSR	PQT	VPG	GP	Eureqa
Nguyen-1	$x^3 + x^2 + x$	100	100	100	96	100	100
Nguyen-2	$x^4 + x^3 + x^2 + x$	100	100	99	47	97	100
Nguyen-3	$x^5 + x^4 + x^3 + x^2 + x$	100	100	86	4	100	95
Nguyen-4	$x^6 + x^5 + x^4 + x^3 + x^2 + x$	100	100	93	1	100	70
Nguyen-5	$\sin(x^2) \cos(x) - 1$	100	72	73	5	45	73
Nguyen-6	$\sin(x) + \sin(x + x^2)$	100	100	98	100	91	100
Nguyen-7	$\log(x + 1) + \log(x^2 + 1)$	97	35	41	3	0	85
Nguyen-8	\sqrt{x}	100	96	21	5	5	0
Nguyen-9	$\sin(x) + \sin(y^2)$	100	100	100	100	100	100
Nguyen-10	$2 \sin(x) \cos(y)$	100	100	91	99	76	64
Nguyen-11	x^y	100	100	100	100	7	100
Nguyen-12	$x^4 - x^3 + \frac{1}{2}y^2 - y$	0	0	0	0	0	0
Average		91.4	83.6	75.2	46.7	60.1	73.9

From *Symbolic Regression via Neural-Guided Genetic Programming Population Seeding* [Neurips 2021]

Take Away

Take Away

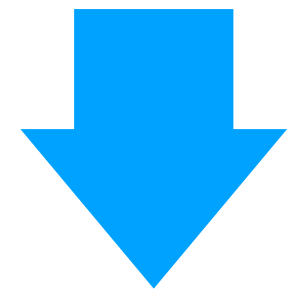
PL problems



Learning problems

Take Away

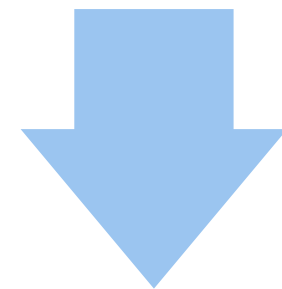
PL problems



Learning problems

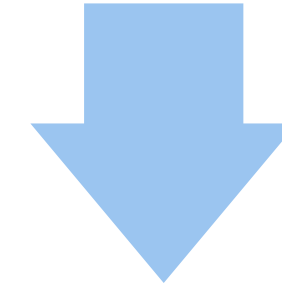
Ex. programs, pre/post-conditions, expectations

A certain kind of maps



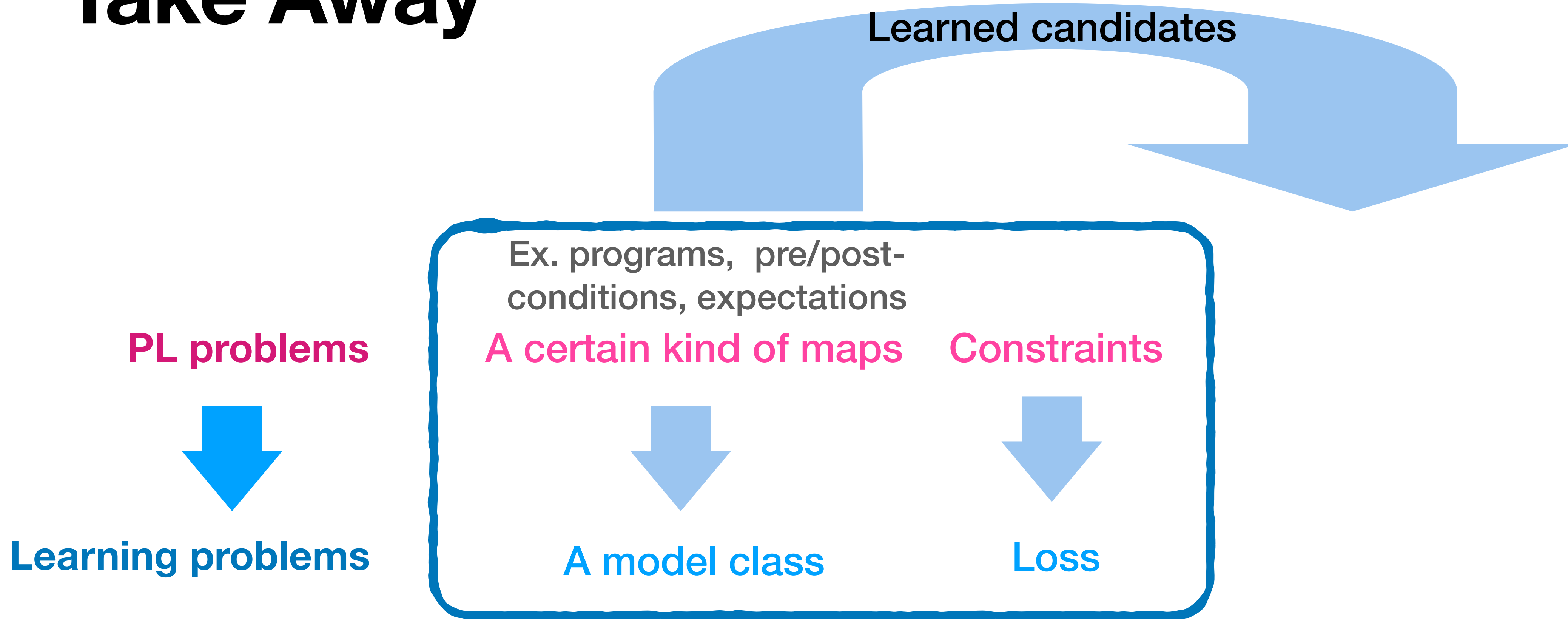
A model class

Constraints



Loss

Take Away



Take Away

